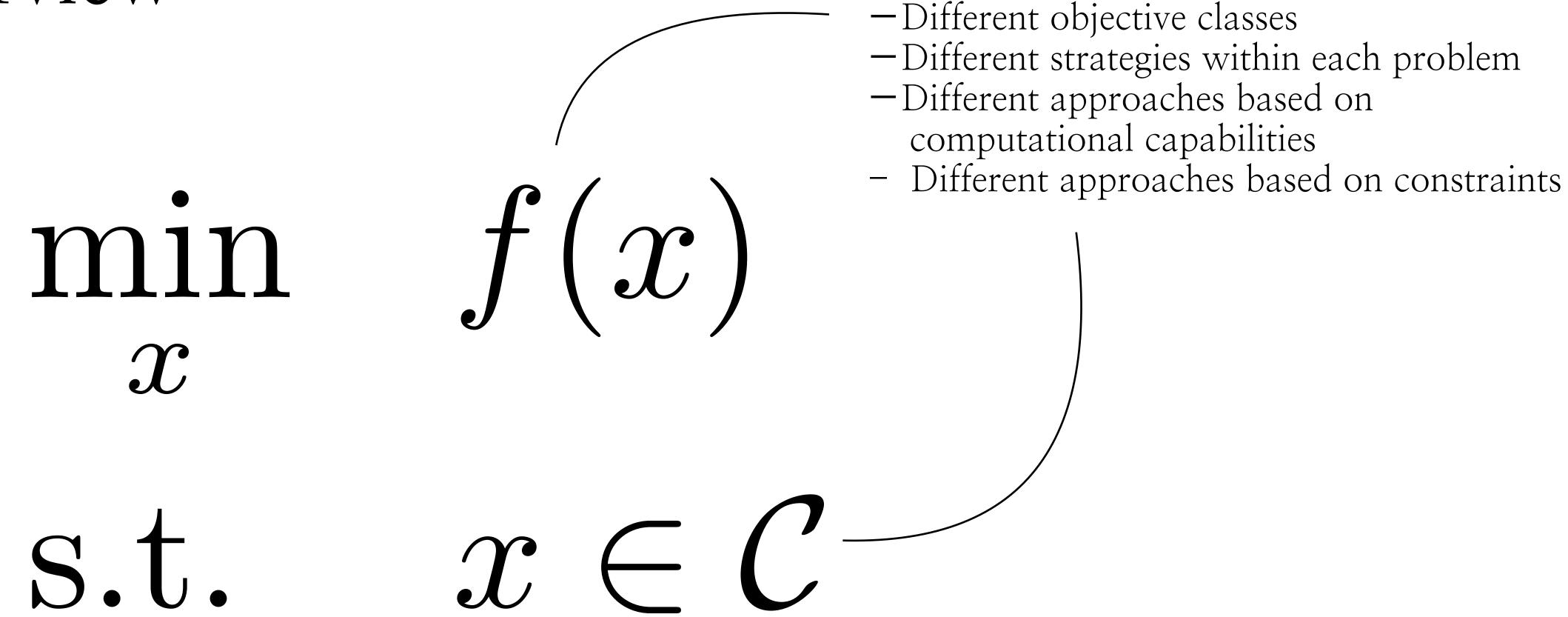
# COMP 414/514: Optimization – Algorithms, Complexity and Approximations

#### Overview



And, always having in mind applications in machine learning, AI and signal processing

s.t.  $x \in C$ 

$$\min_{x} f(x)$$

Unconstrained optimization

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#### Overview

- In this lecture, we will:
  - Discuss how to distribute optimization in large-scale settings
  - Study synchrony vs. asynchrony in gradient descent
  - Provide some rough theoretical results on how asynchrony affects performance
  - Alternatives and state of the art

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- Properties: i) the current model  $x_t$  is used for the computation of  $\nabla f_{i_t}(\cdot)$  ii) when we update the model, the state of the system is as when we read  $x_t$  iii) The whole process is sequential

$$x_{t+1} = x_t - \eta \nabla f_{i_t}(x_t) = x_{t-1} - \eta \left( \nabla f_{i_t}(x_t) + \nabla f_{i_{t-1}}(x_{t-1}) \right) = \dots = x_0 - \eta \sum_{i} \nabla f_{i_j}(x_j)$$

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- How can we distribute this computation over multiple processing units?

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- Single node distributed computing:
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- Multi-node distributed computing:
  - i) Many machines (up to 1000s), probably with many cores each
     ii) Shared-nothing architecture (each machine has its own CPU, storage)
     iii) Communication between nodes is much less cheap than single node

- Consider the full gradient descent case:

$$x_{t+1} = x_t - \eta \sum_{i=1}^{N} \nabla f_i(x_t)$$

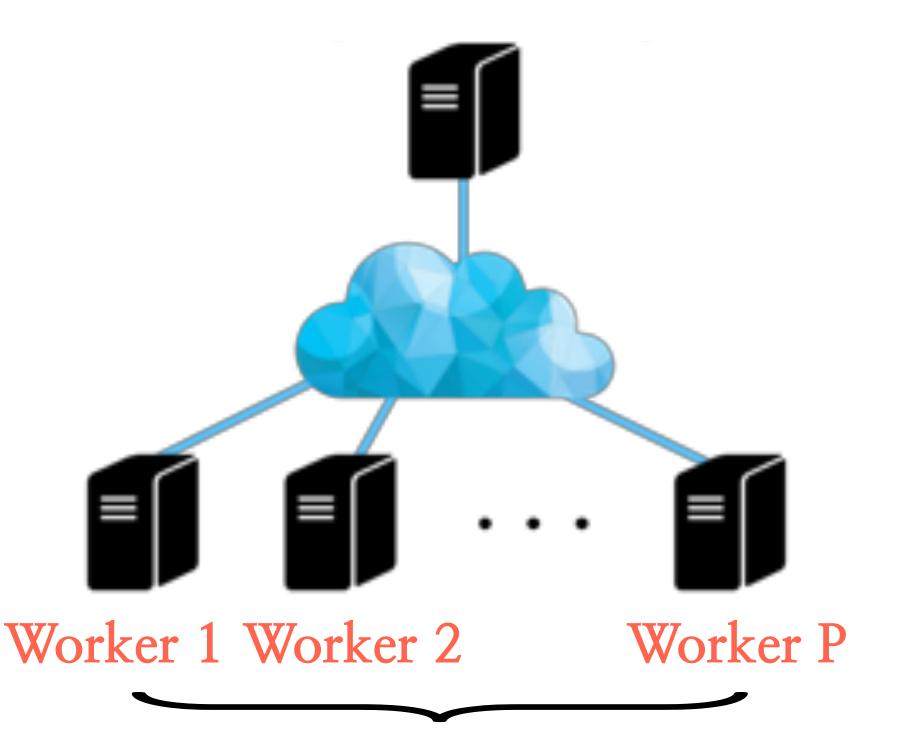
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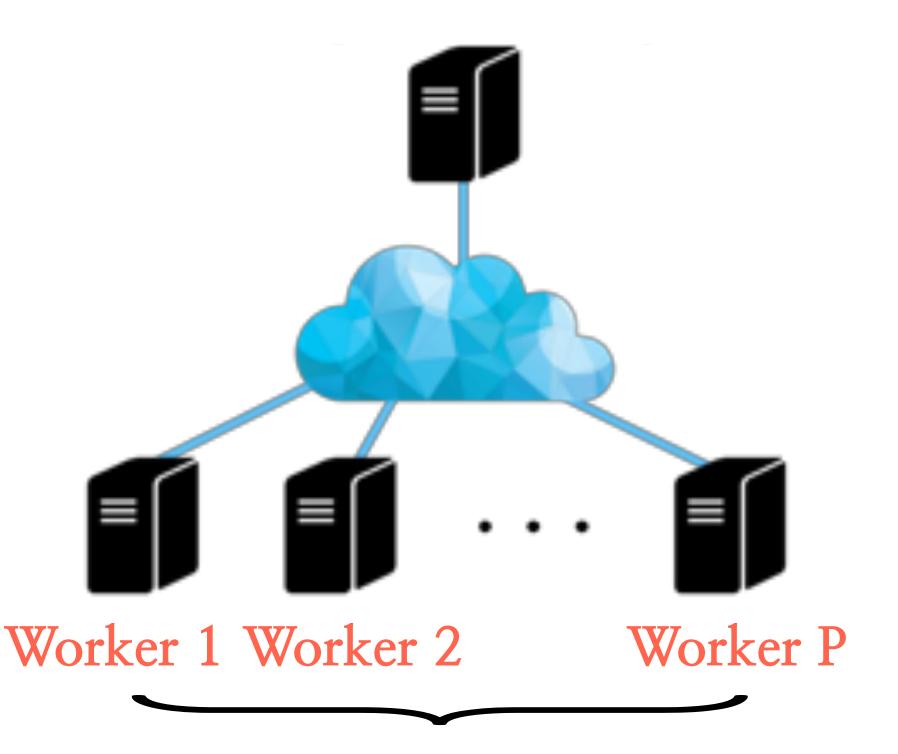


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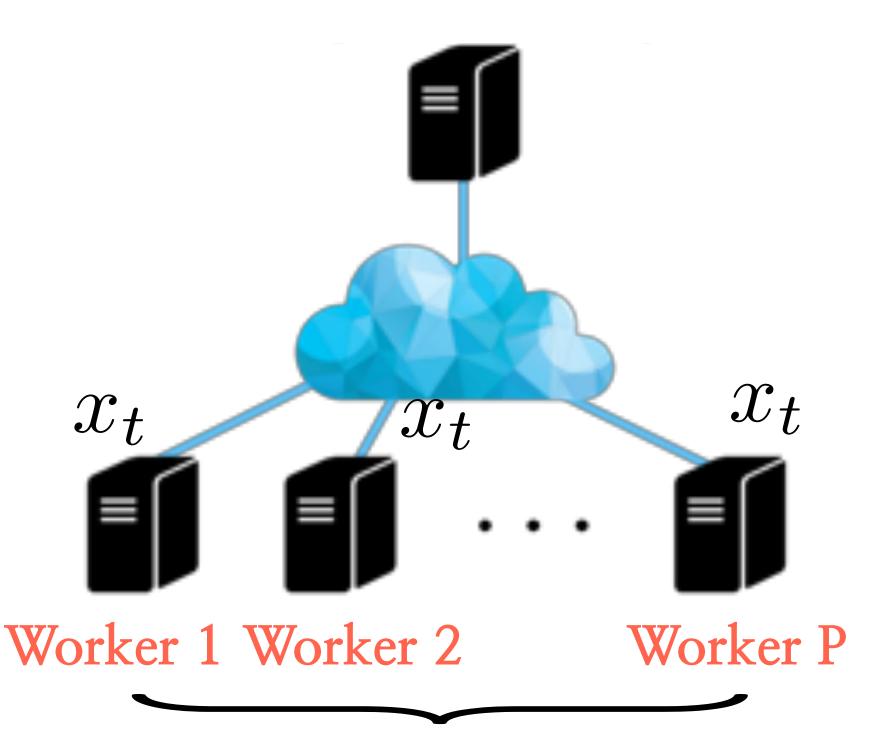
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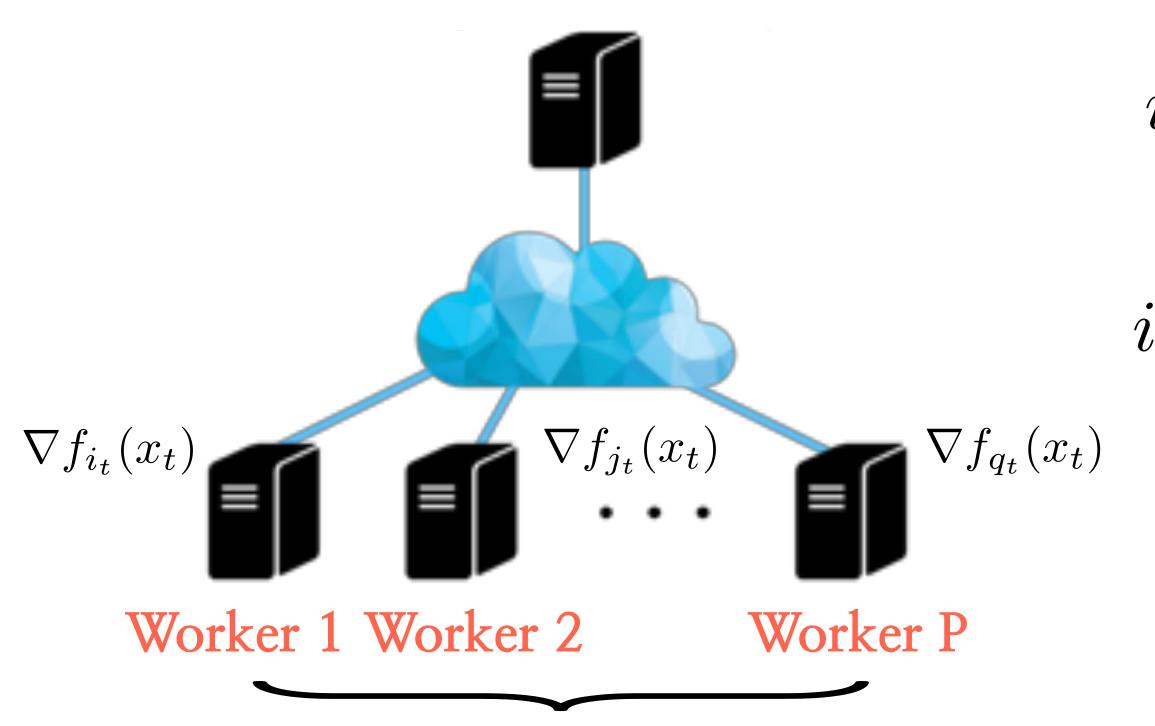
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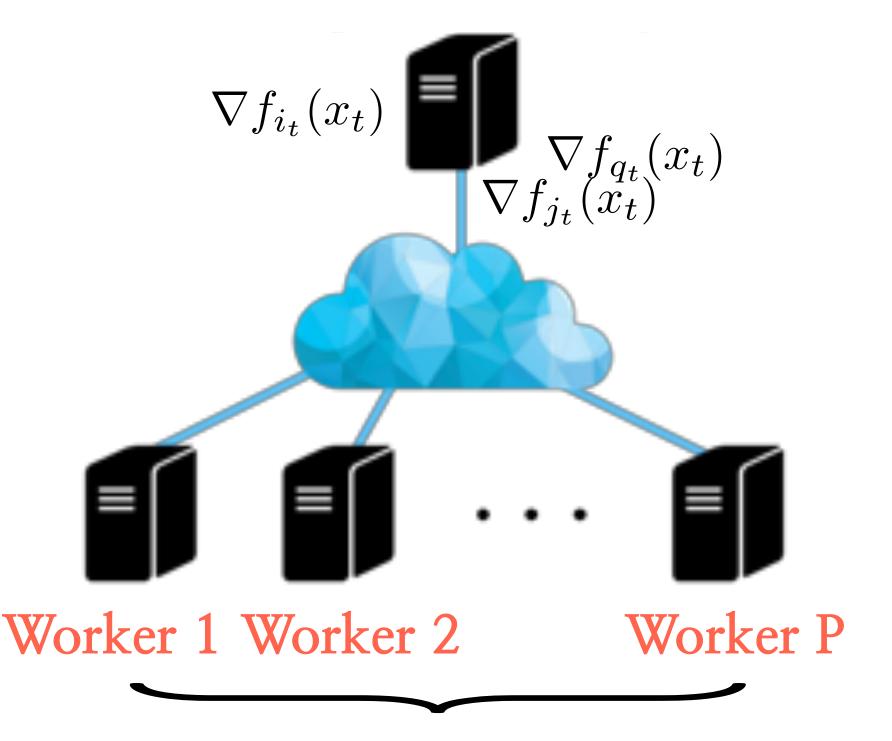
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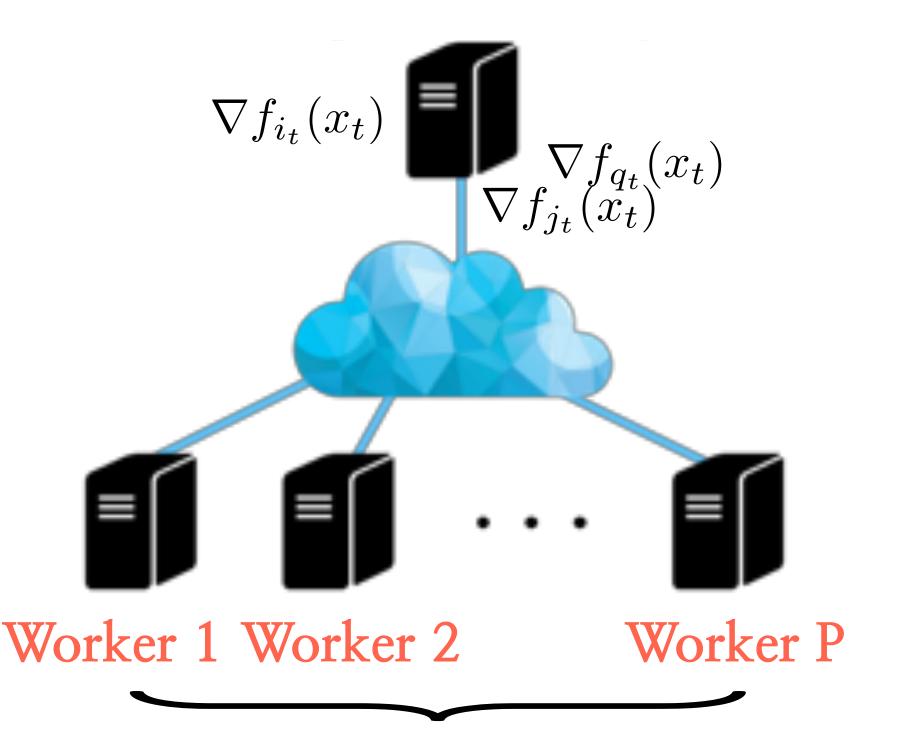
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(..till the very last slow worker – active research: tackle stranglers)

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- "But, there are cases where we have finite & fixed data see neural networks"
- "Well, the problem here is that full gradient descent does not perform well"

### Generalization vs. training error:

- 1. If we care about only the training error, full GD could work well
- 2. In ML tasks, we often care about the generalization error, i.e., the performance of the model on unseen data

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Whiteboard

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### Whiteboard

(This relates to the question 'large vs. small batch training')

- Consider the case where even  $\nabla f_{i_t}(x_t) \in \mathbb{R}^p$  is expensive for a single node

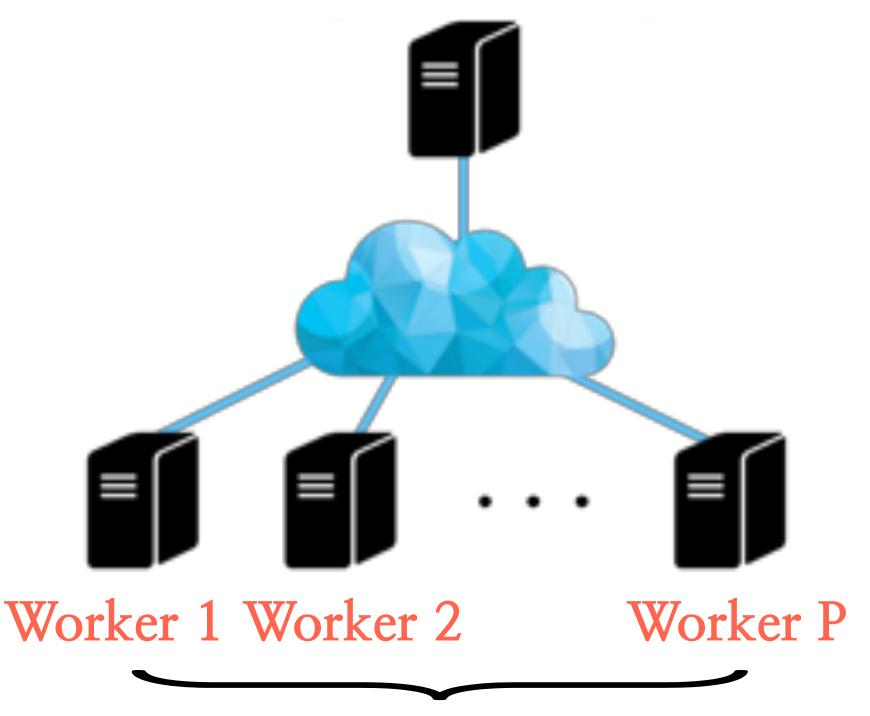
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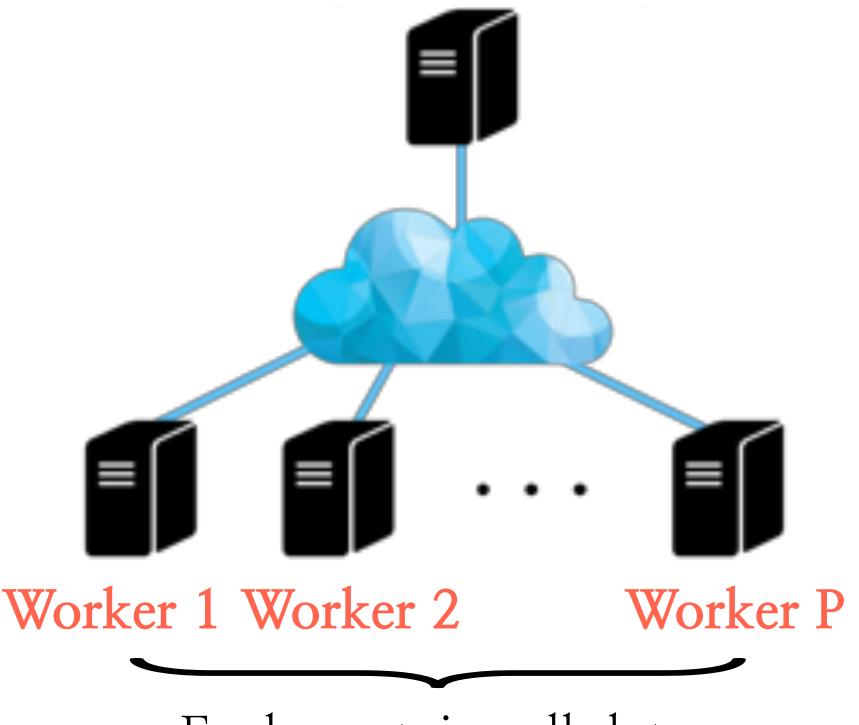


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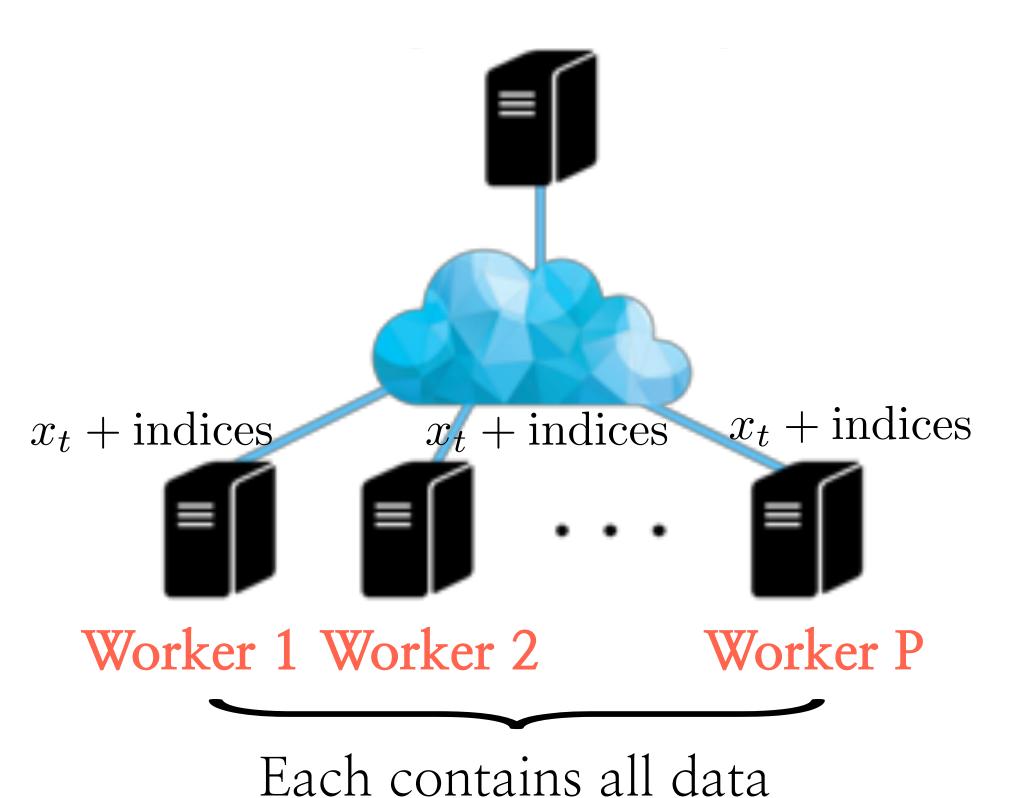
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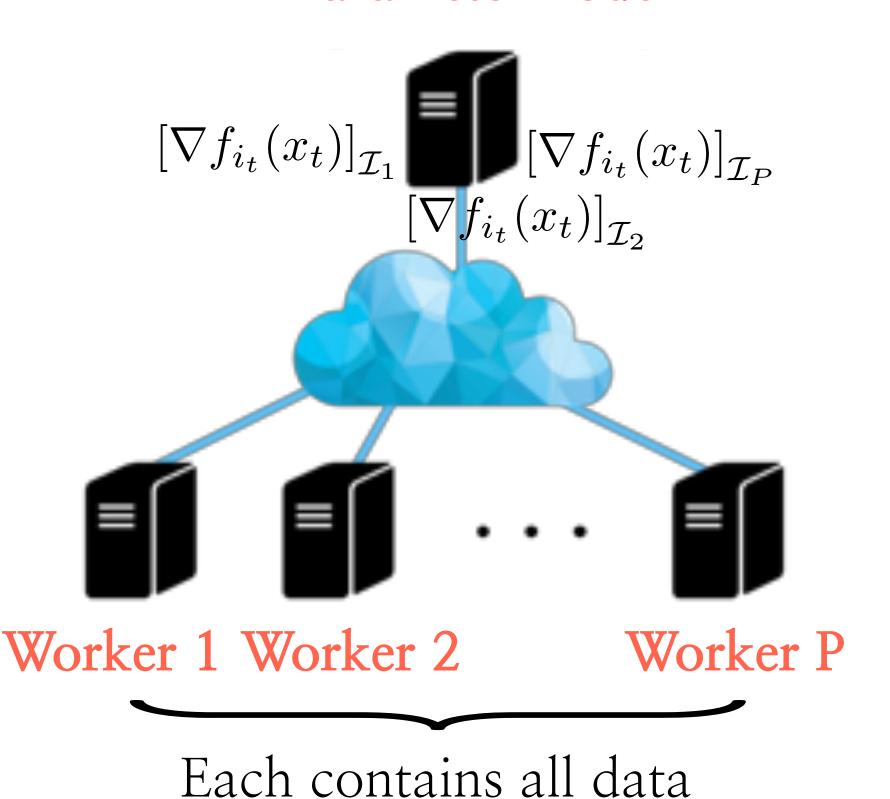
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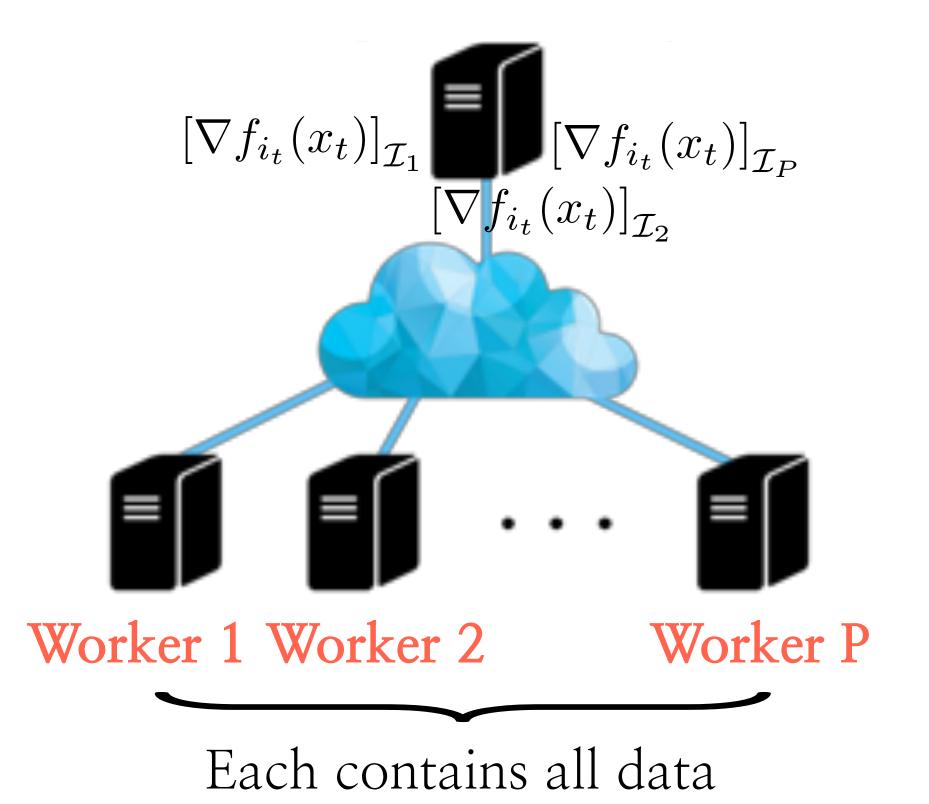


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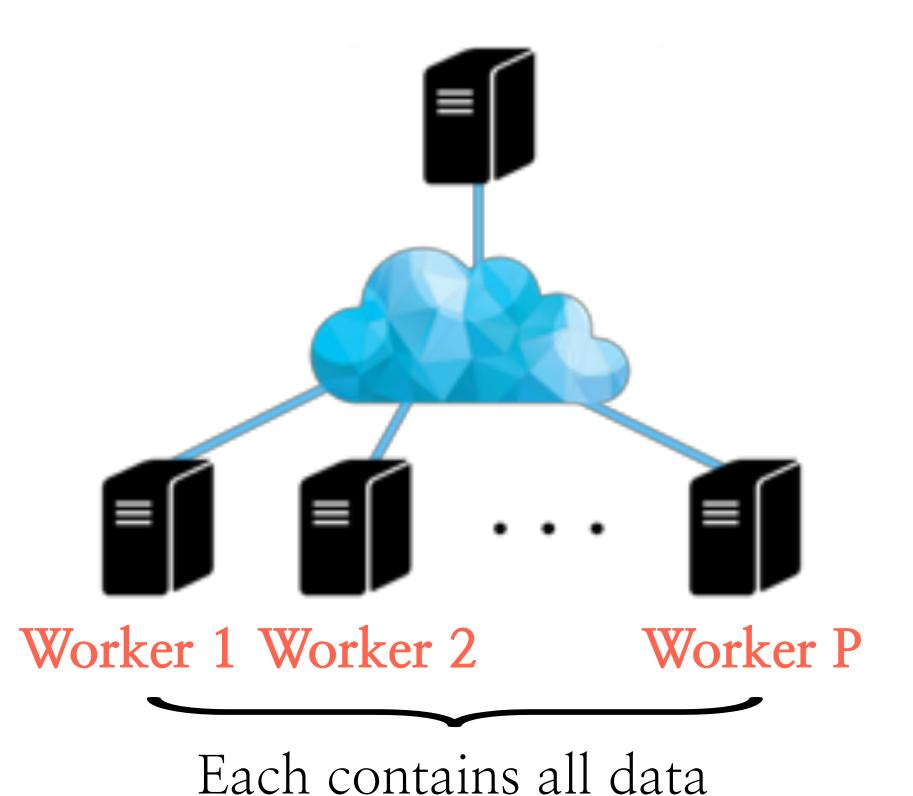
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(..till the very last slow worker)

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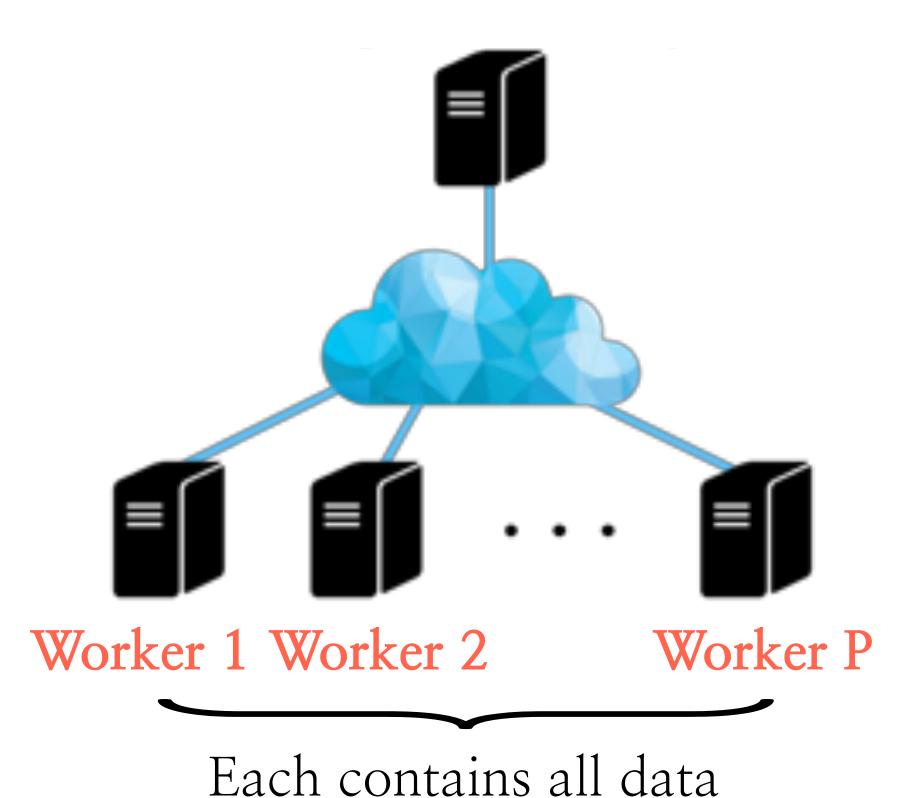
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i) Relates to coordinate descent algorithms

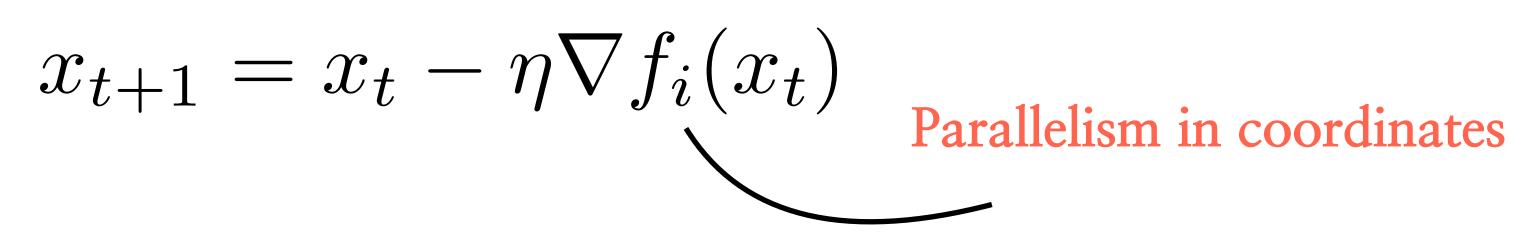
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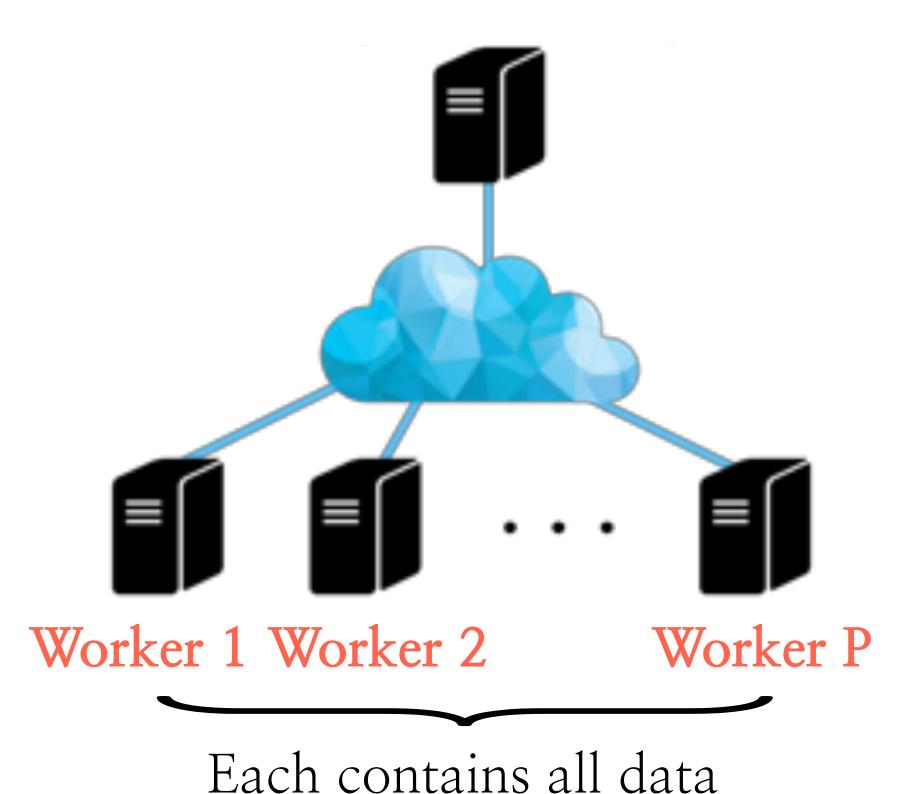
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- i) Relates to coordinate descent algorithms
- *ii*) Could be part of a large-scale implementation, where part of the model is too large to be computed in a centralized fashion
- iii) Could be an overkill to only compute updates for a subset of entries

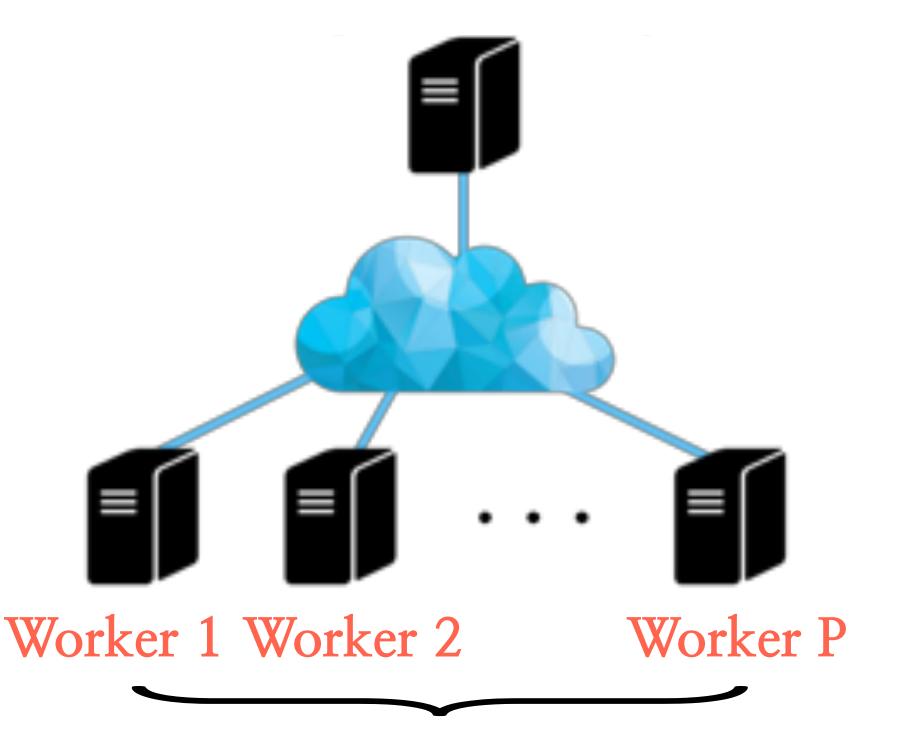
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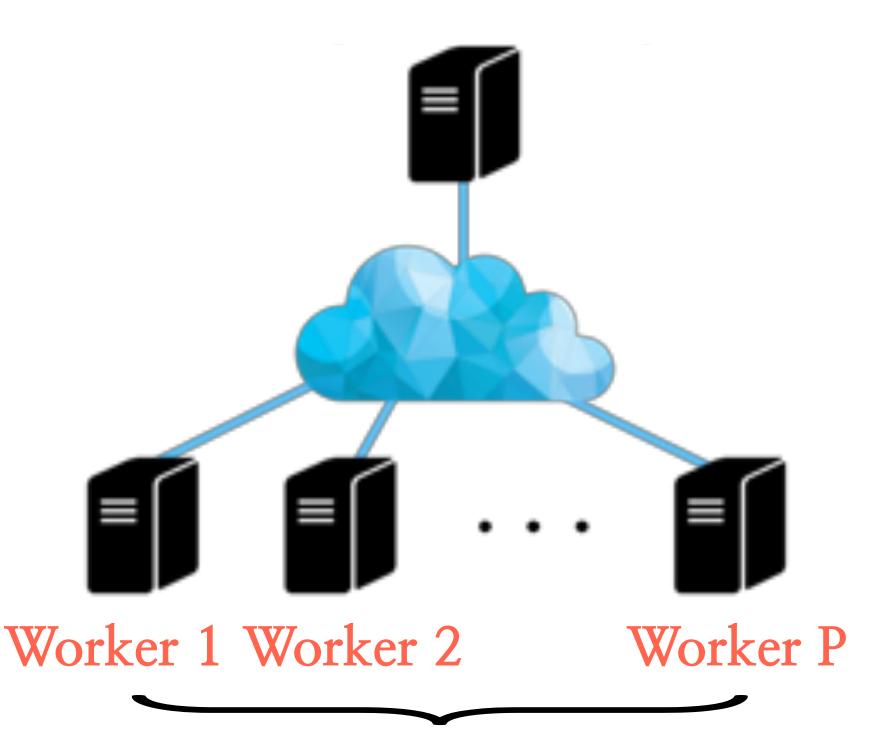


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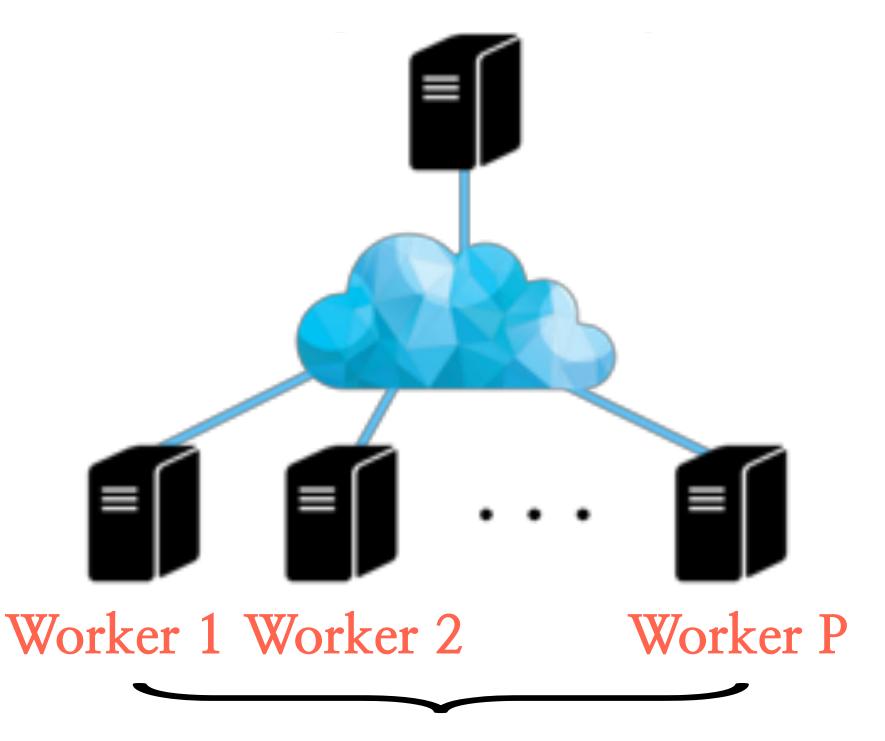


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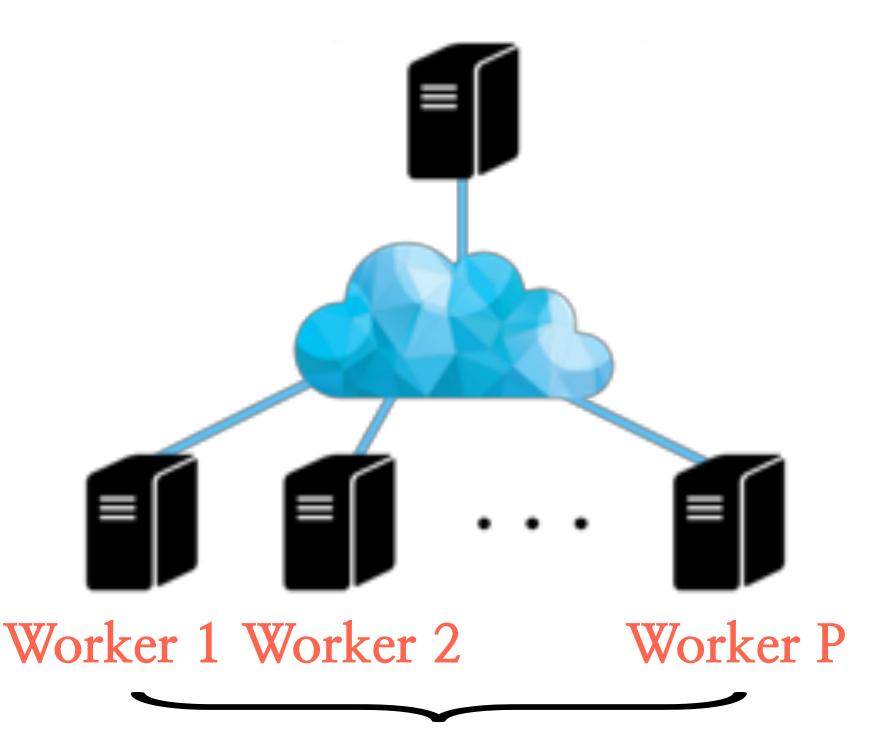


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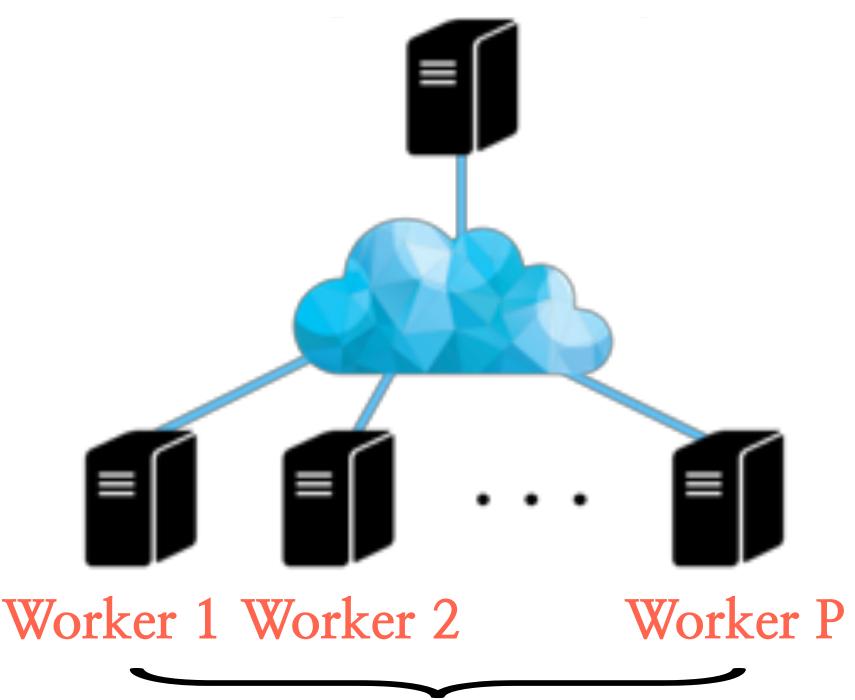
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- What about the setting in-between? Mini-batch SGD

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i) Still requires synchronization; each worker has less work to do

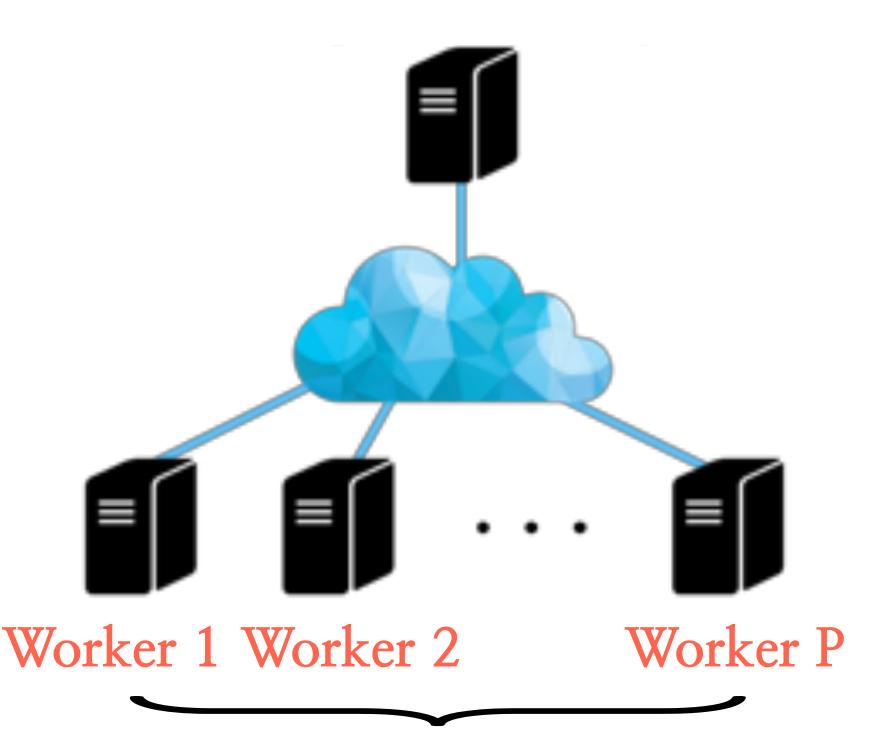
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*ii*) Introduces a tradeoff between statistical efficiency, computations efficiency (in terms of convergence) and communication efficiency

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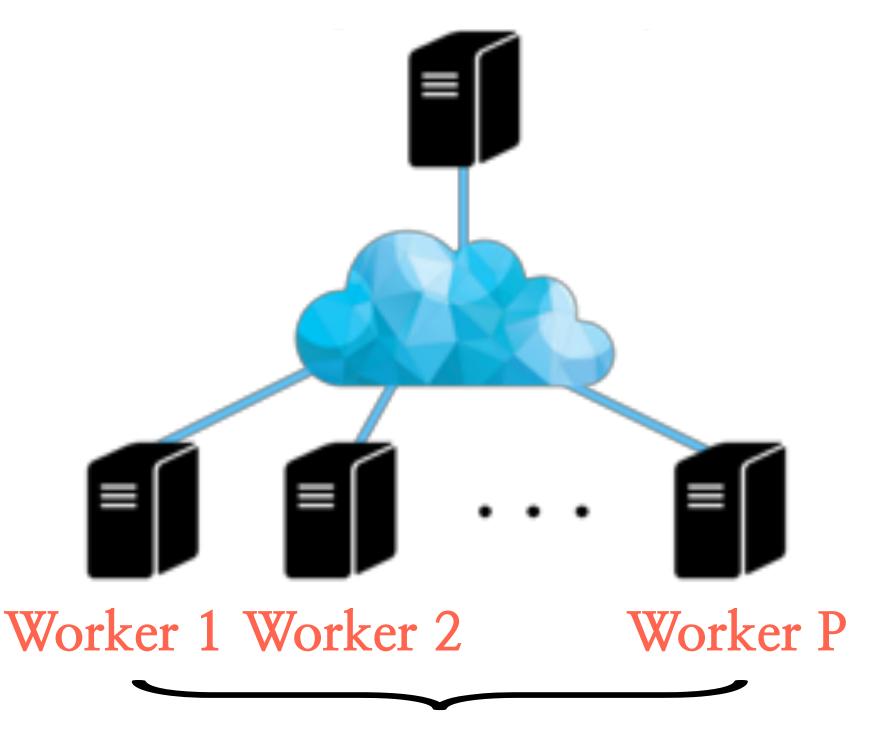
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- i) Still requires synchronization; each worker has less work to do
- ii) Introduces a tradeoff between statistical efficiency, computations efficiency (in terms of convergence) and communication efficiency
- iii) Usually computing  $\nabla f_{i_t}(x_t)$  is cheap per node

(Discussion about large batch training)

- What if we run mini-batch SGD in parallel and combine at the end:

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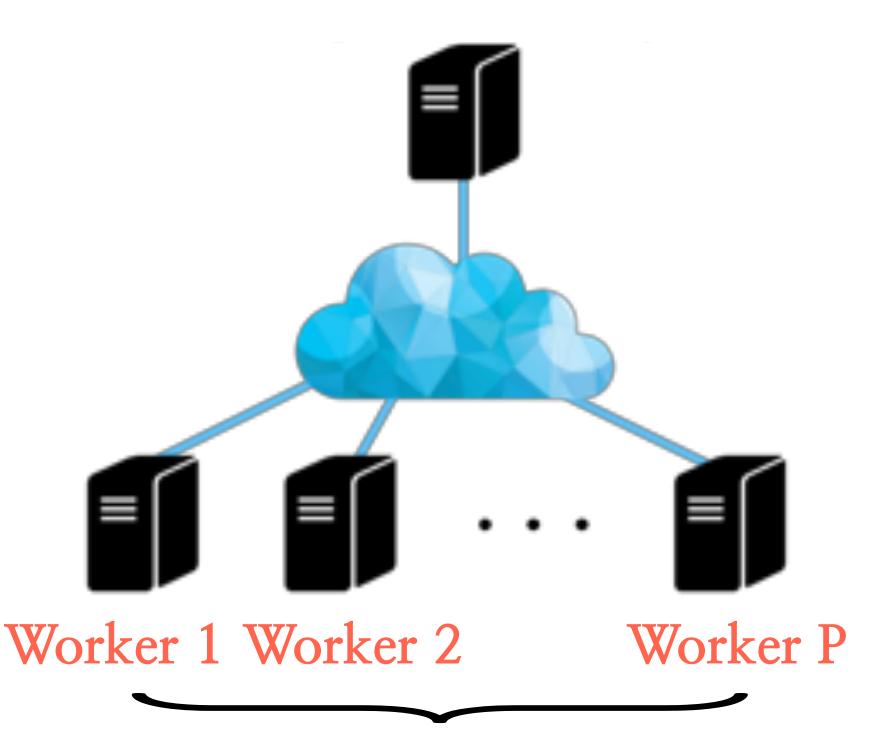
Worker 1 Worker 2 Worker P

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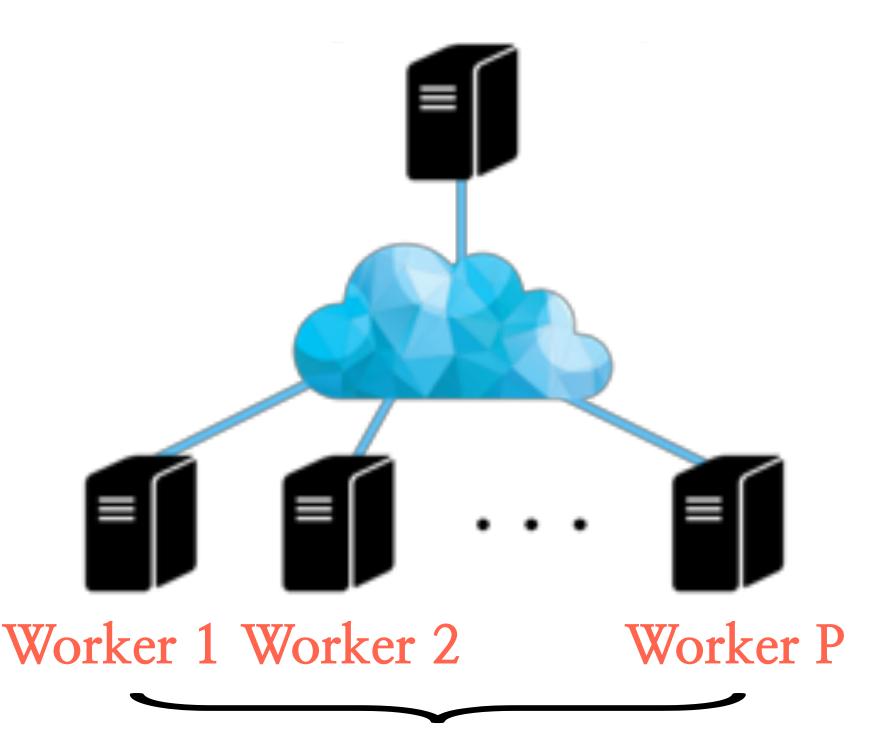


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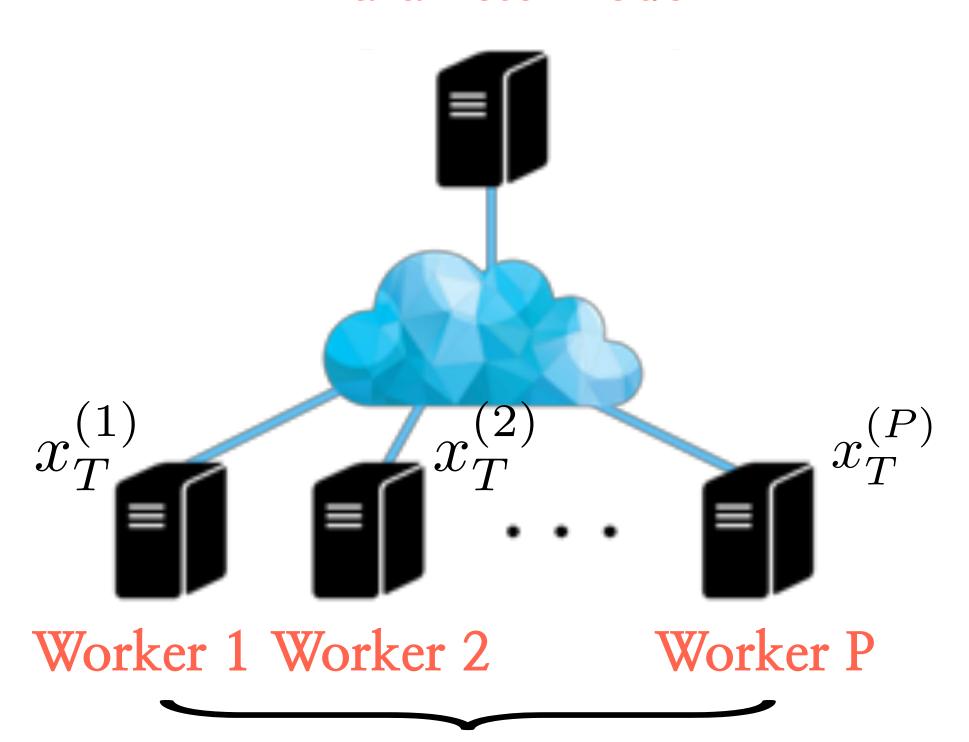
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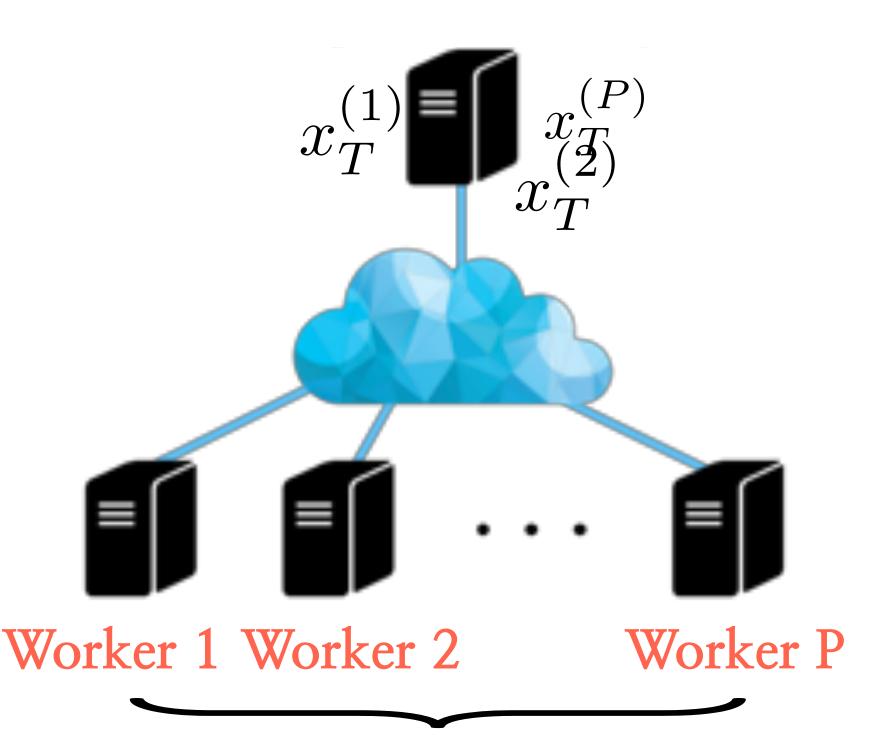
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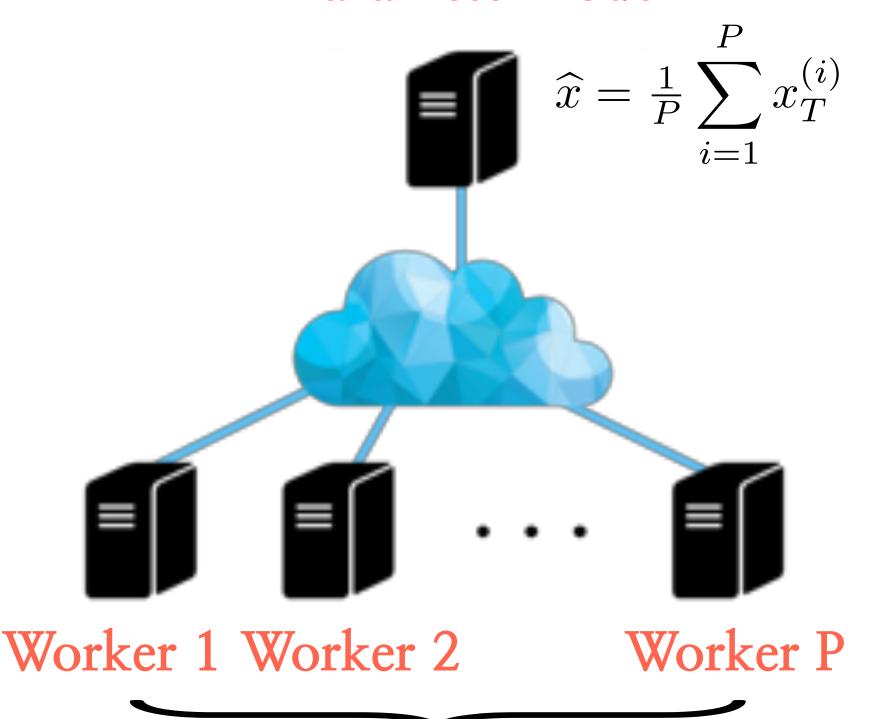
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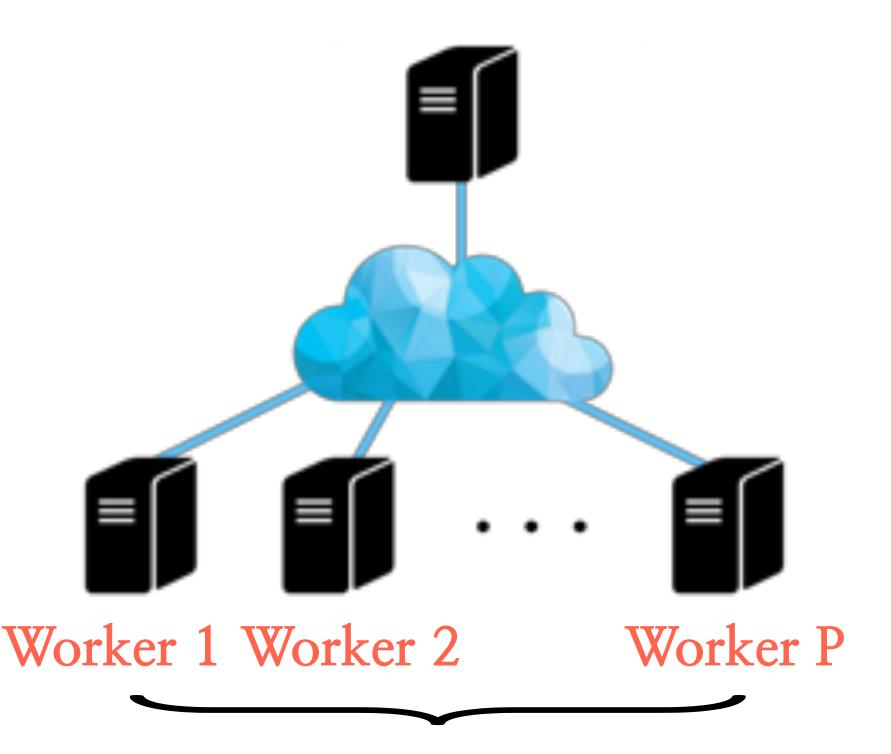
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- What if we run mini-batch SGD in parallel and combine at the end:

$$x_{t+1} = x_t - \eta \sum_{i \in \mathcal{I}_t} \nabla f_i(x_t)$$

#### Parameter node

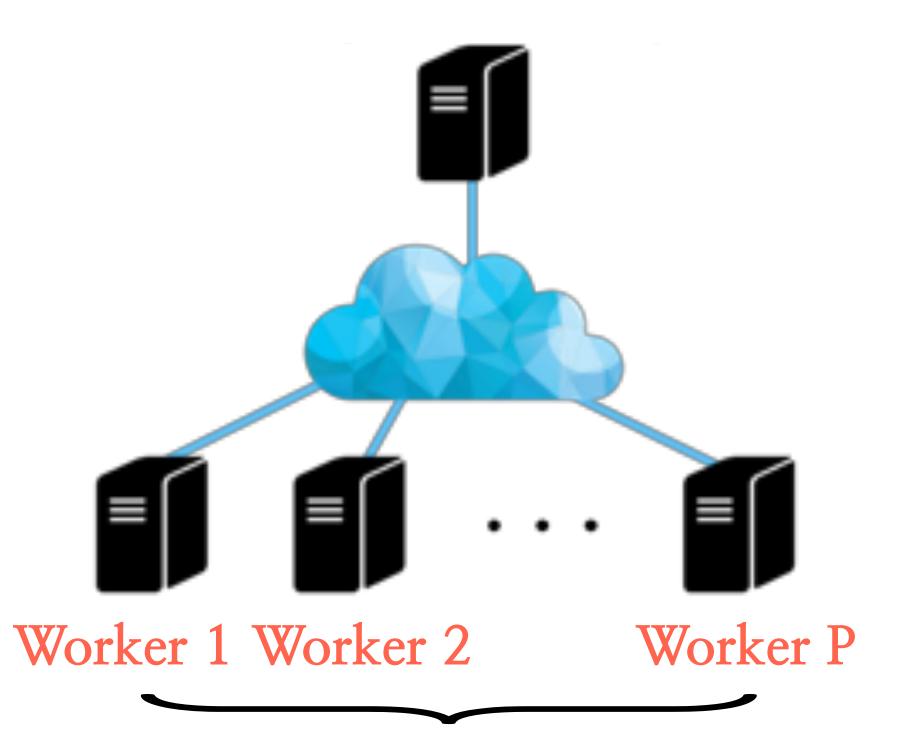


Each contains distinct partition of data

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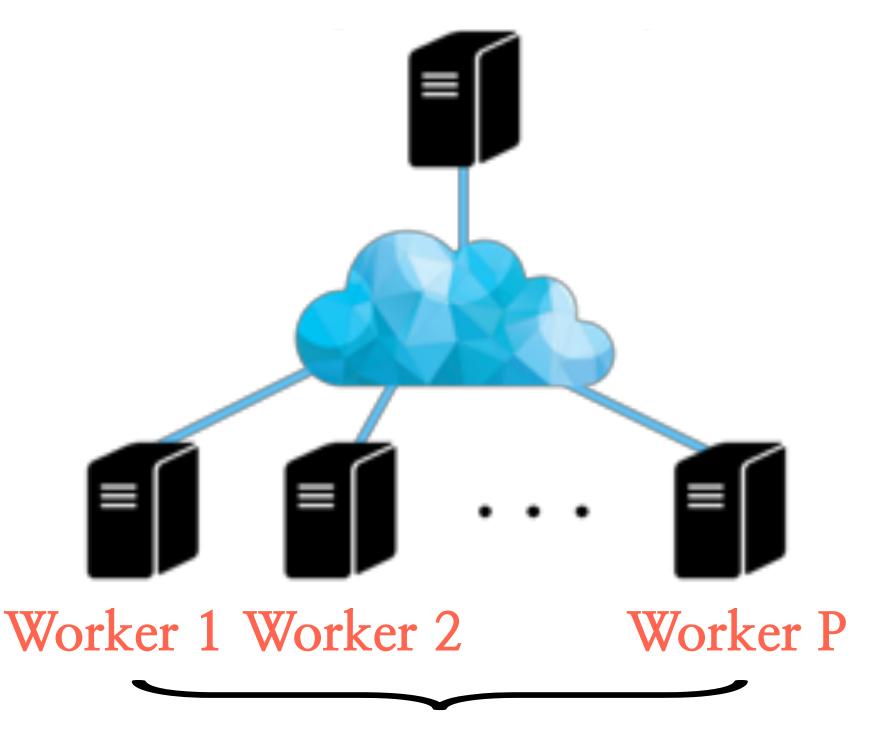


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- ii) The model was designed for convex problems the idea is that each subproblem has a solution close to the global one thus averaging does not hurt
- iii) Final decision is prediction averaging similar ideas hold for random forests

# Using distributed computing in a different way

- Run code in parallel as a way for hyperparameter optimization

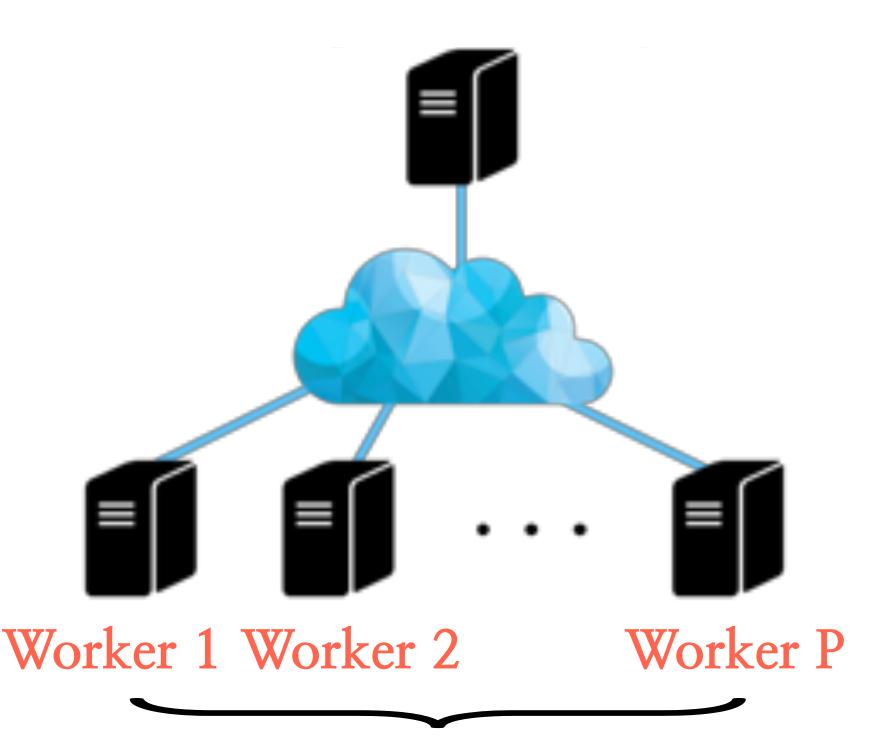
$$x_{t+1} = x_t - \eta_1 \sum_{i \in \mathcal{I}_t} \nabla f_i(x_t) \qquad \bullet \quad \bullet \quad \bullet \qquad x_{t+1} = x_t - \eta_q \sum_{i \in \mathcal{I}_t} \nabla f_i(x_t)$$

# Using distributed computing in a different way

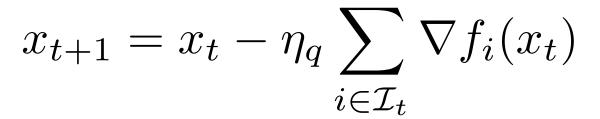
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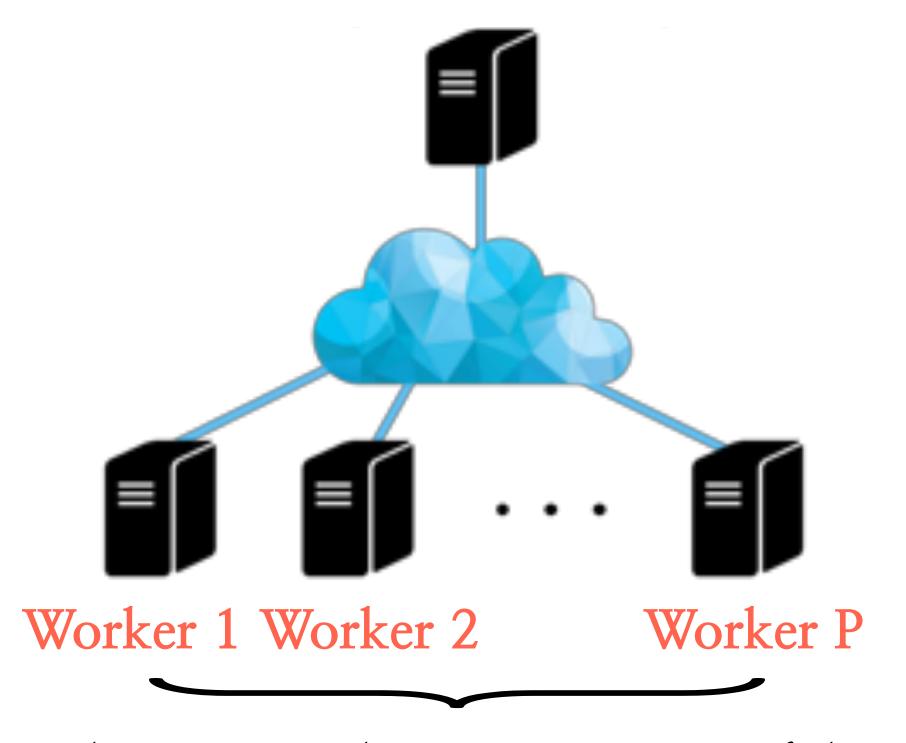
Parameter node



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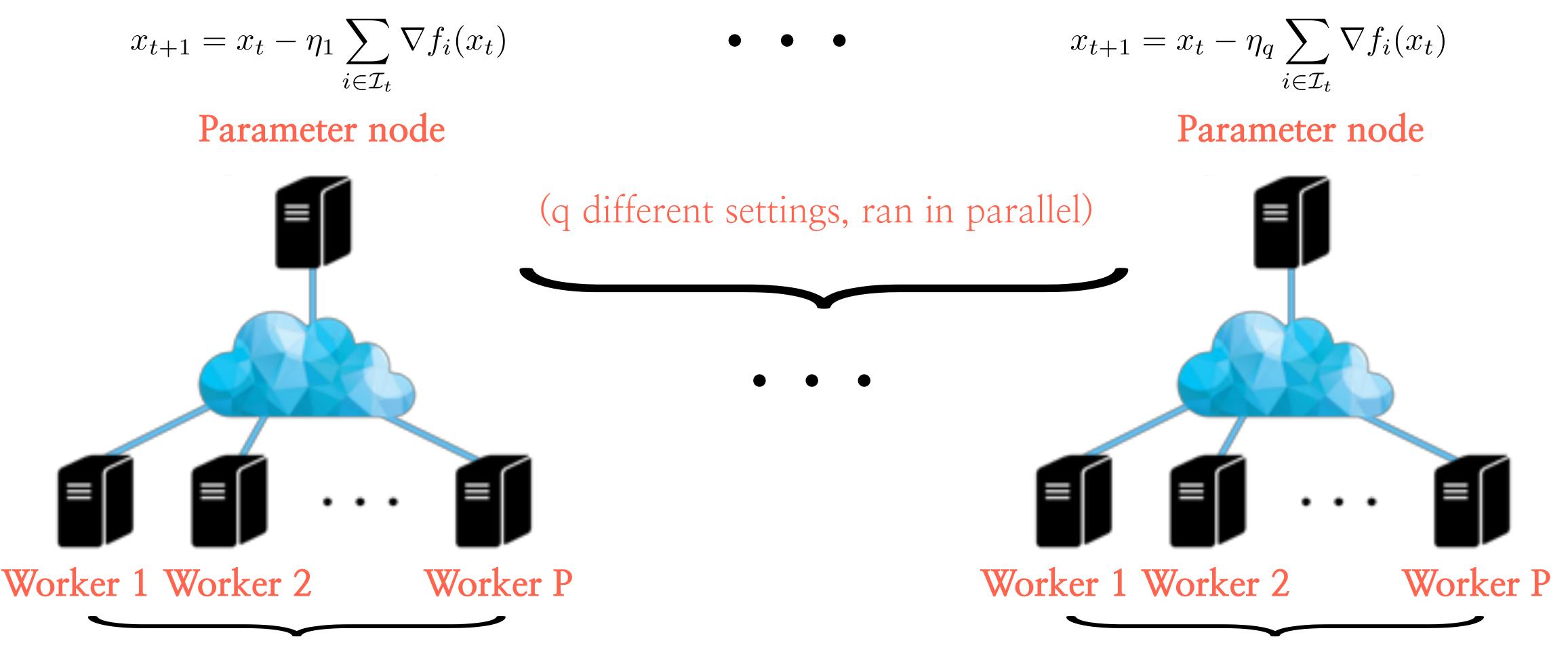
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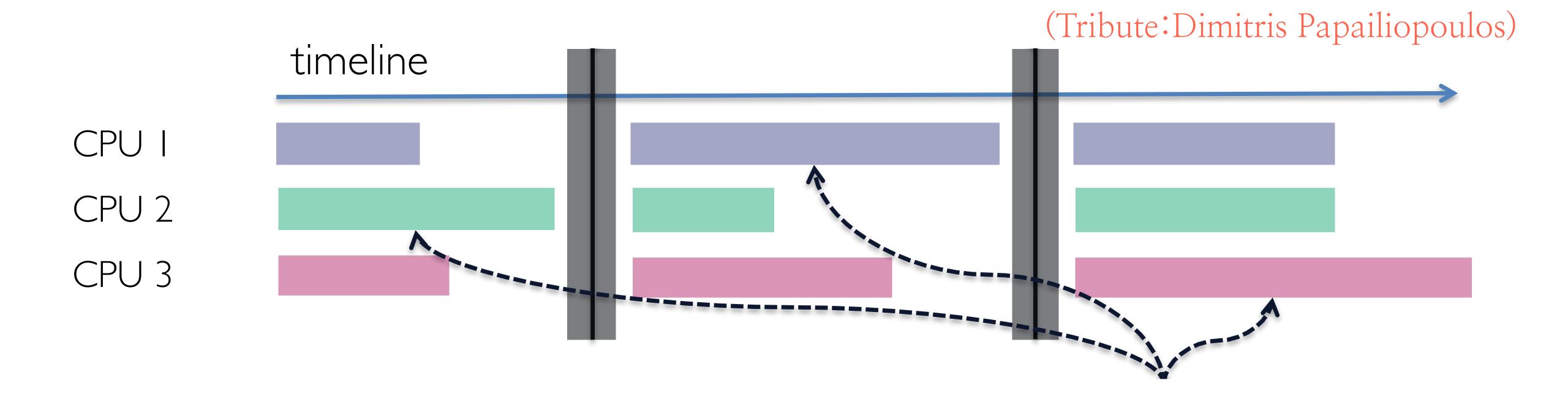
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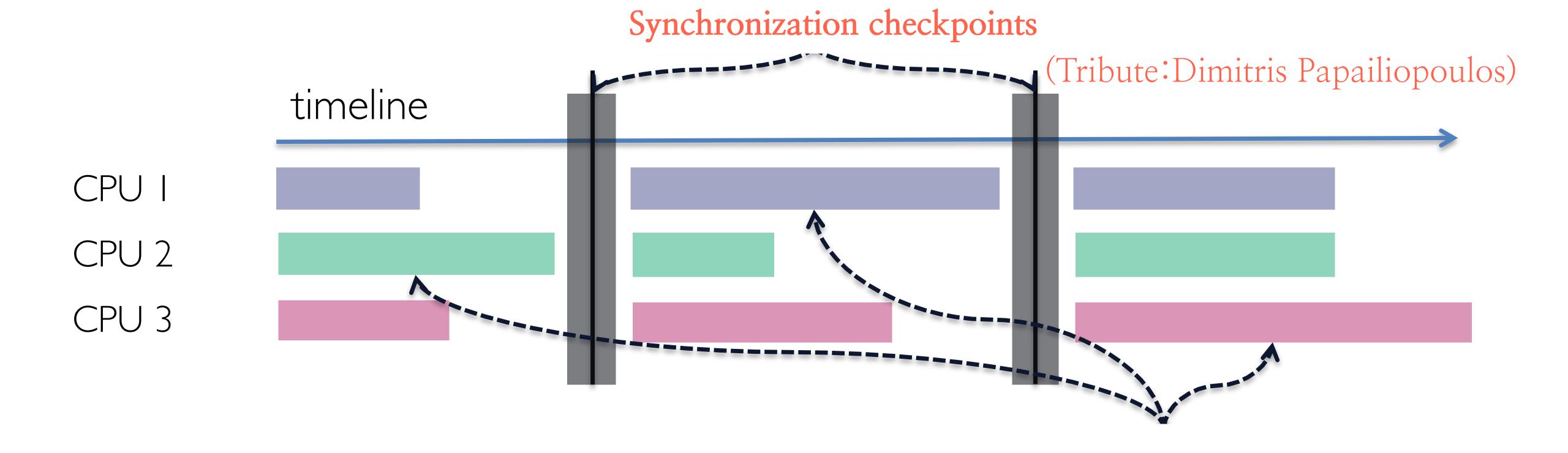
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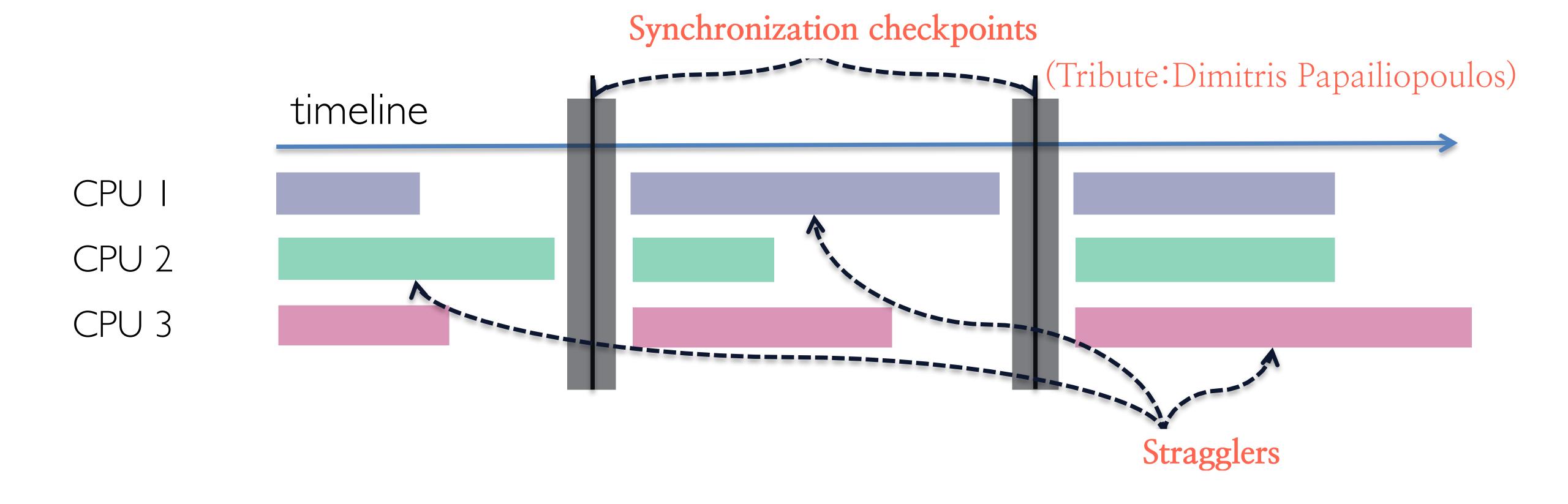


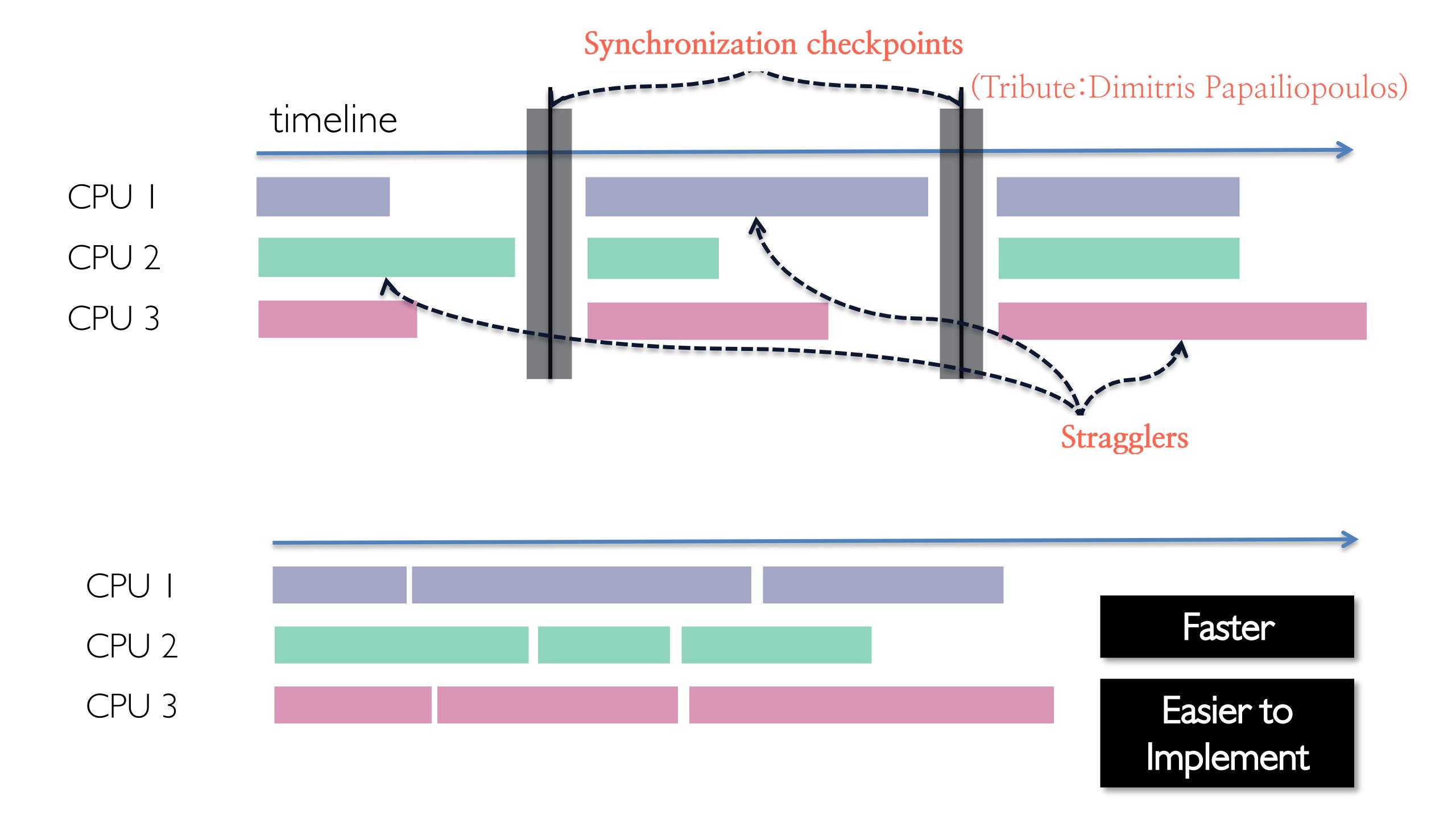
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- Alternatives or we have to bear with this situation?

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#### - Main bottleneck:

- Synchronization (locking) amongst processors

- Run SGD in parallel without locks!





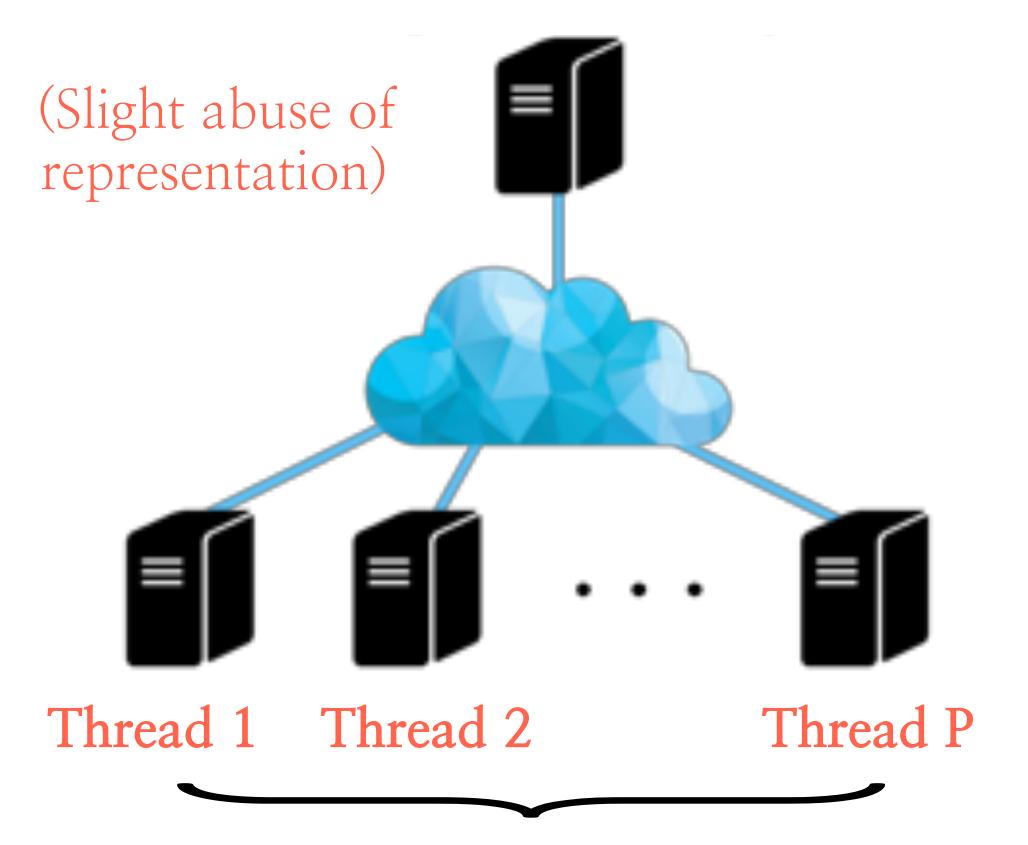


- Run SGD in parallel without locks!



Shared memory

- All threads have access to shared memory

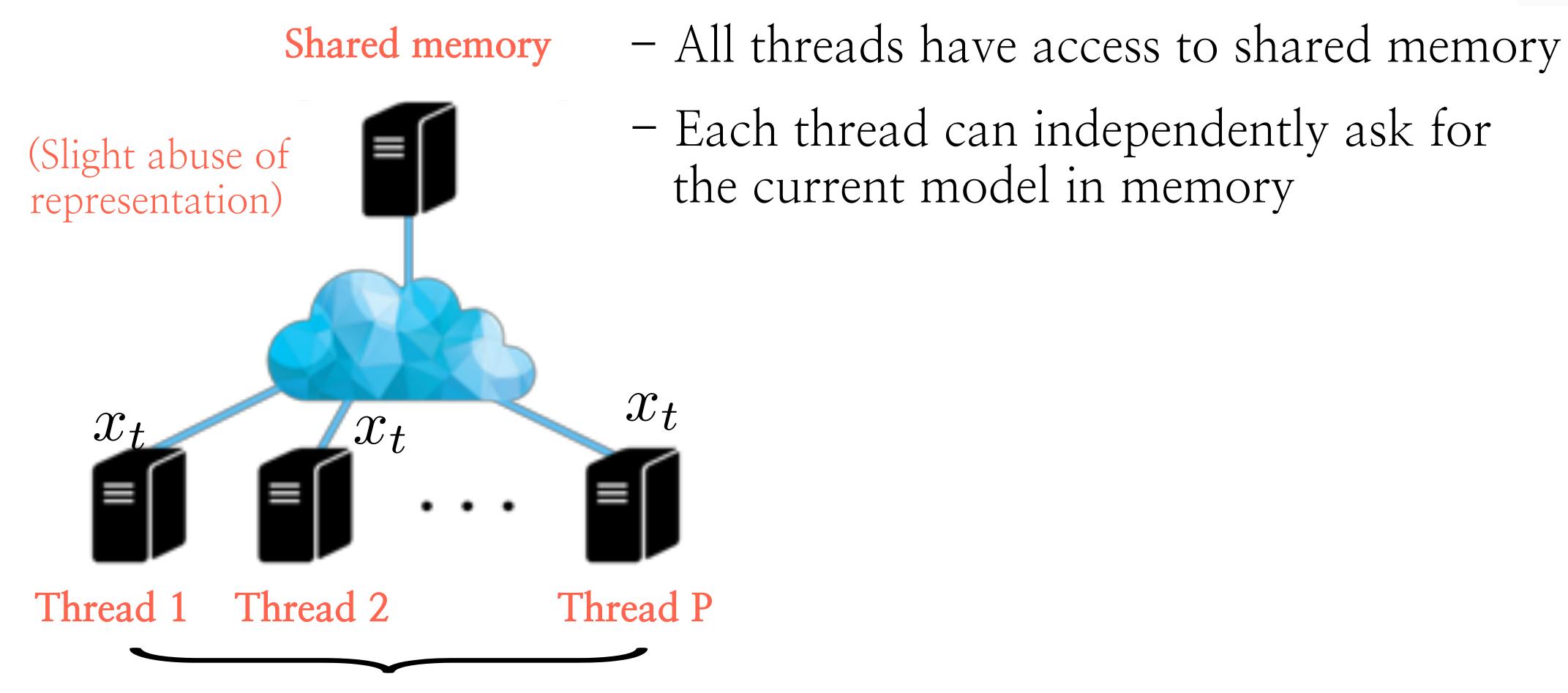


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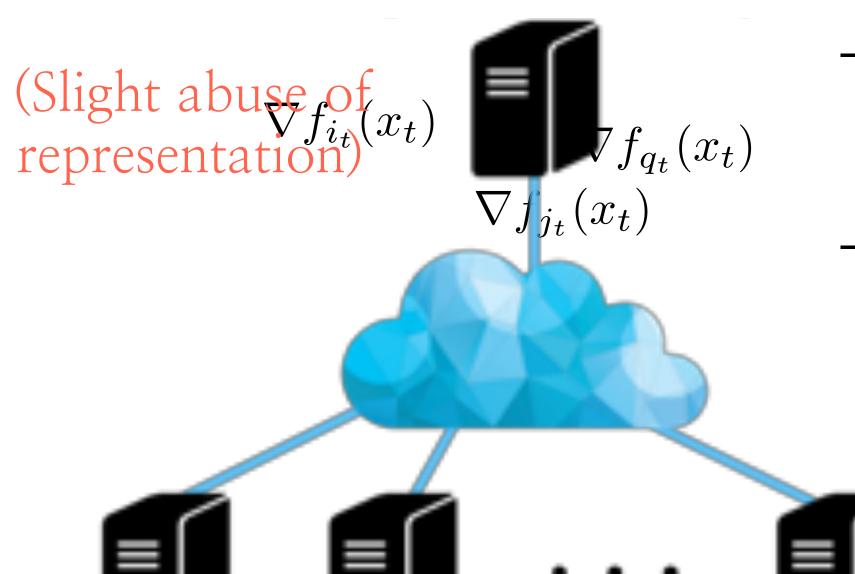
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### - Run SGD in parallel without locks!







Thread 1 Thread 2

- All threads have access to shared memory
- Each thread can independently ask for the current model in memory
- Each thread computes an update (=gradient) and then updates shared memory

(the order that updates are sent is random)

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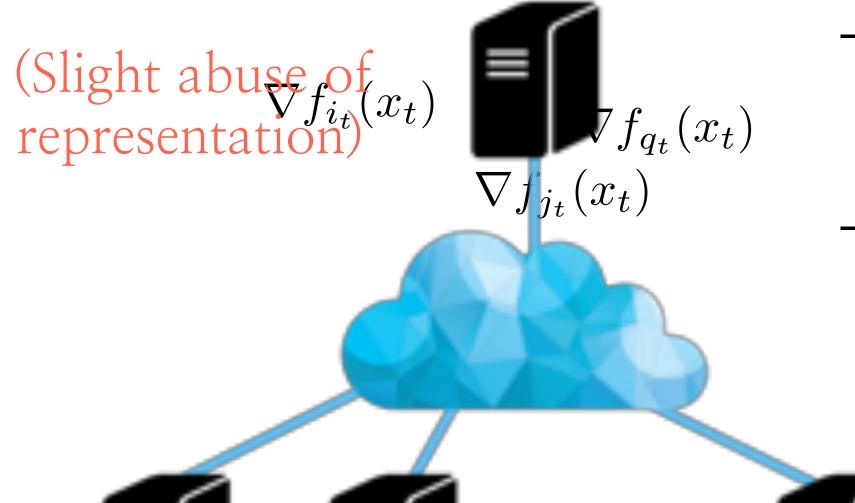
Thread P



#### Microsoft

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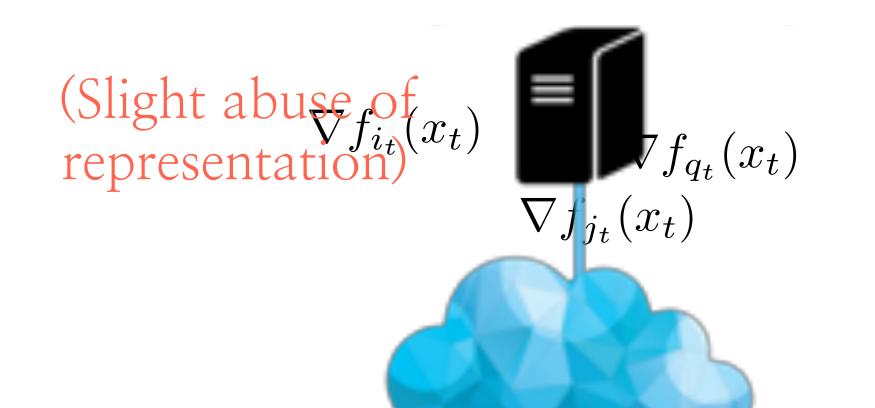
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- The controller in shared memory updates the model in a first-in-first-served fashion
- Assuming all threads have collected  $x_t$

$$x_{t+1} = x_t - \eta \left( \nabla f_{i_t}(x_t) + \nabla f_{j_t}(x_t) + \dots + \nabla f_{q_t}(x_t) \right)$$

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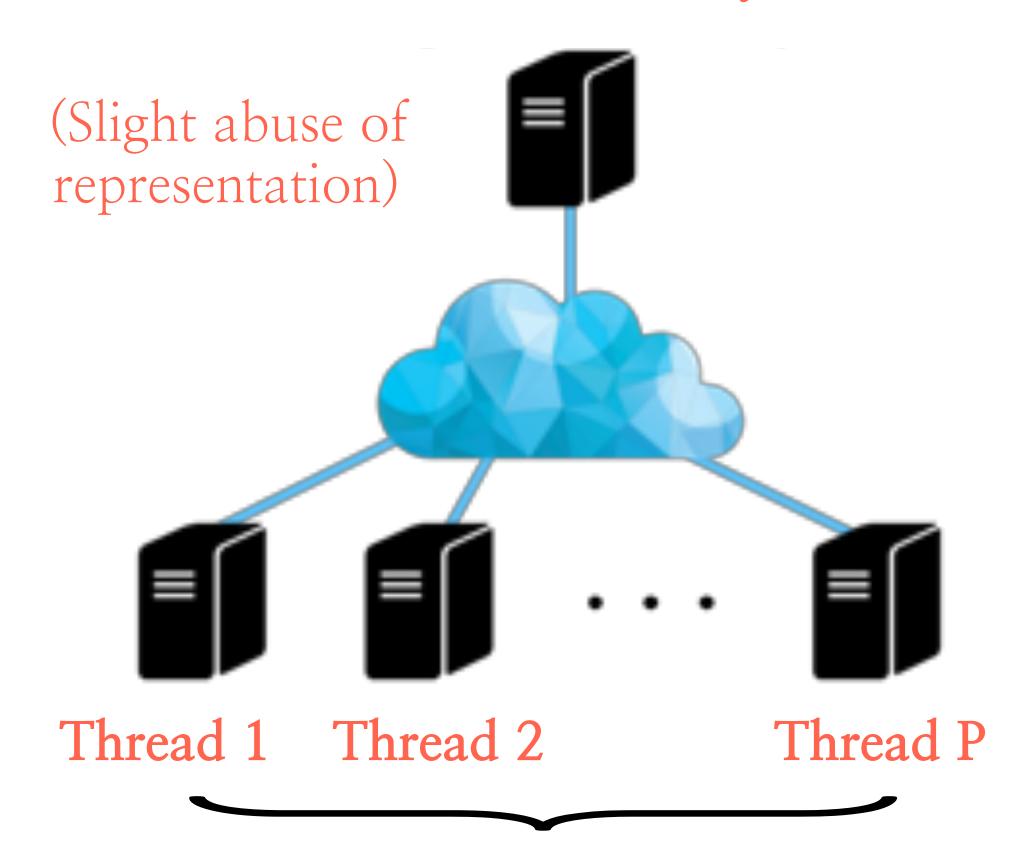
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Google

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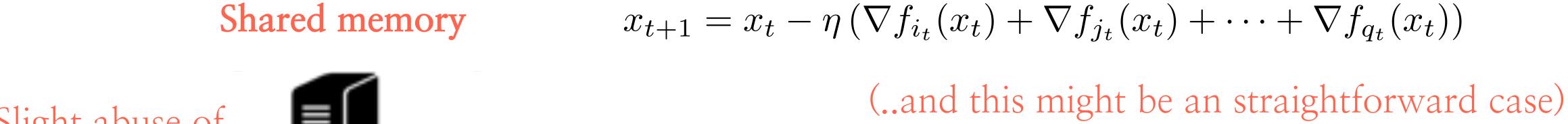
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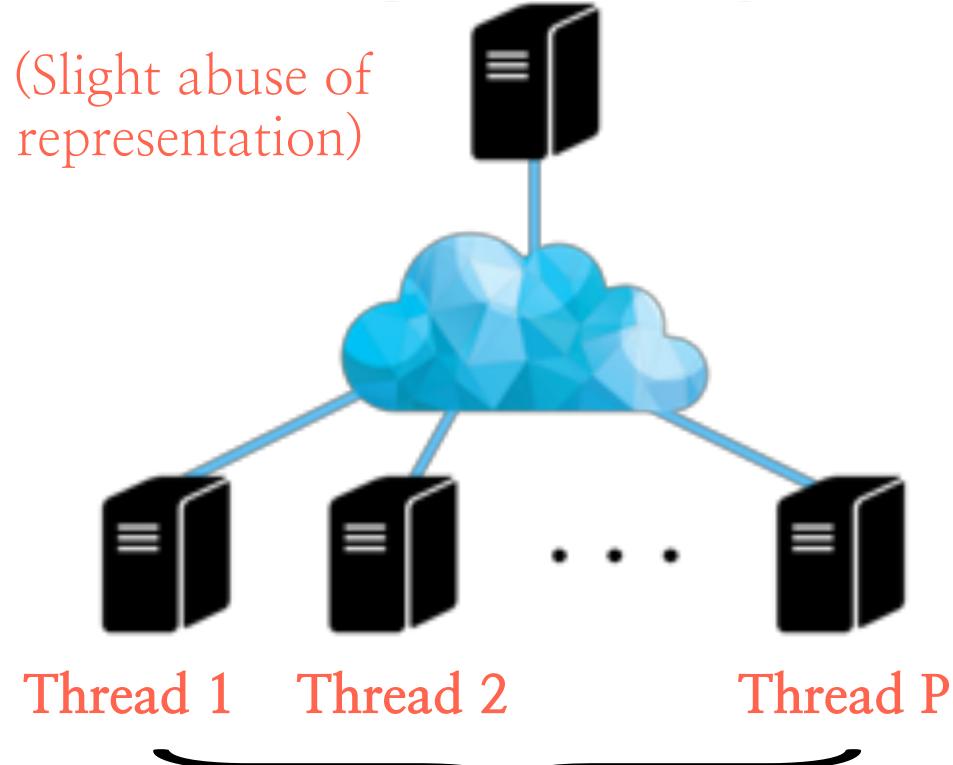


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(Slight abuse of representation) — Three  $x_{t-10}$  —  $x_{t-3}$  —  $x_{t-3}$ 

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Thread P

Thread 1 Thread 2

(..and this might be an straightforward case)

- Threads might process an older model version



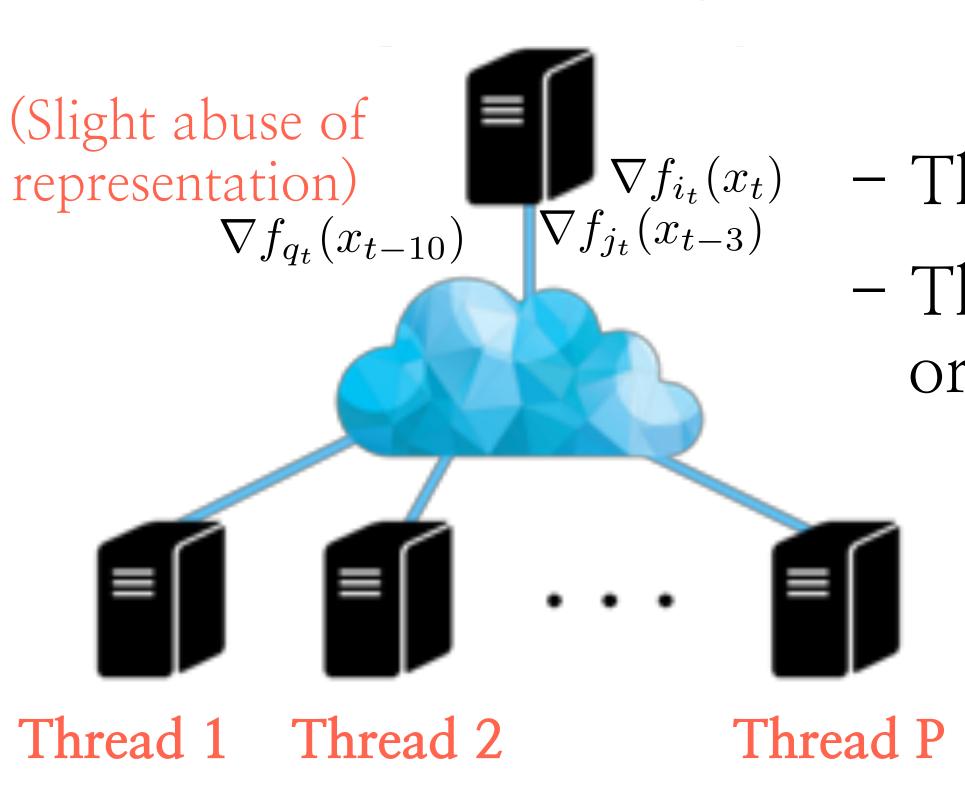
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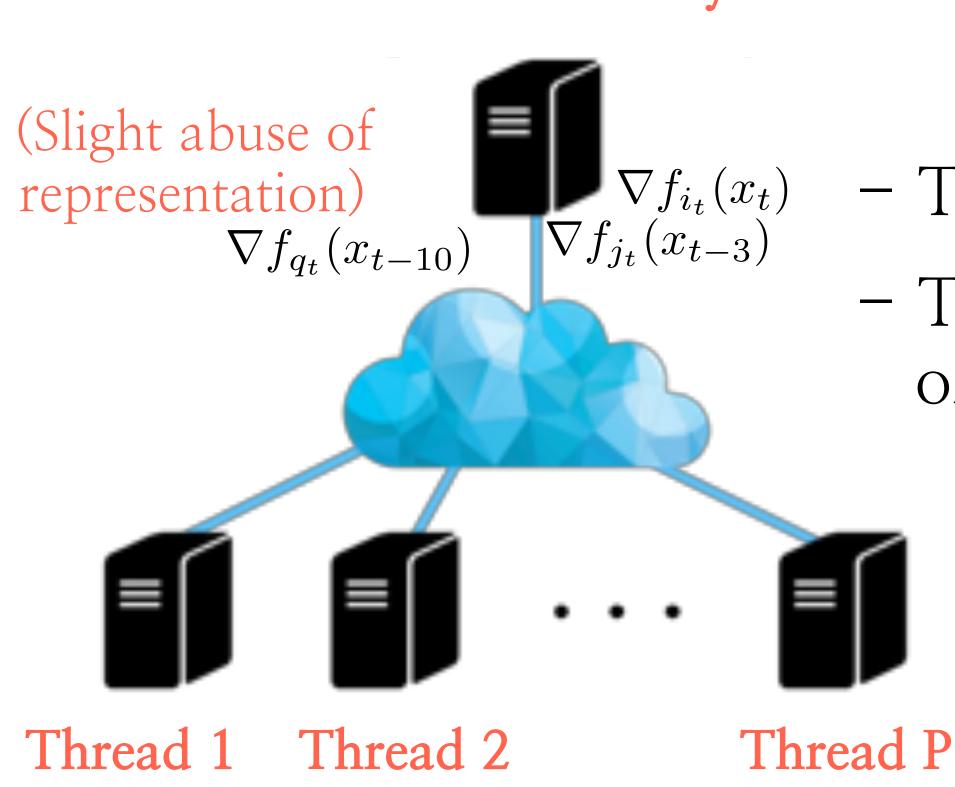
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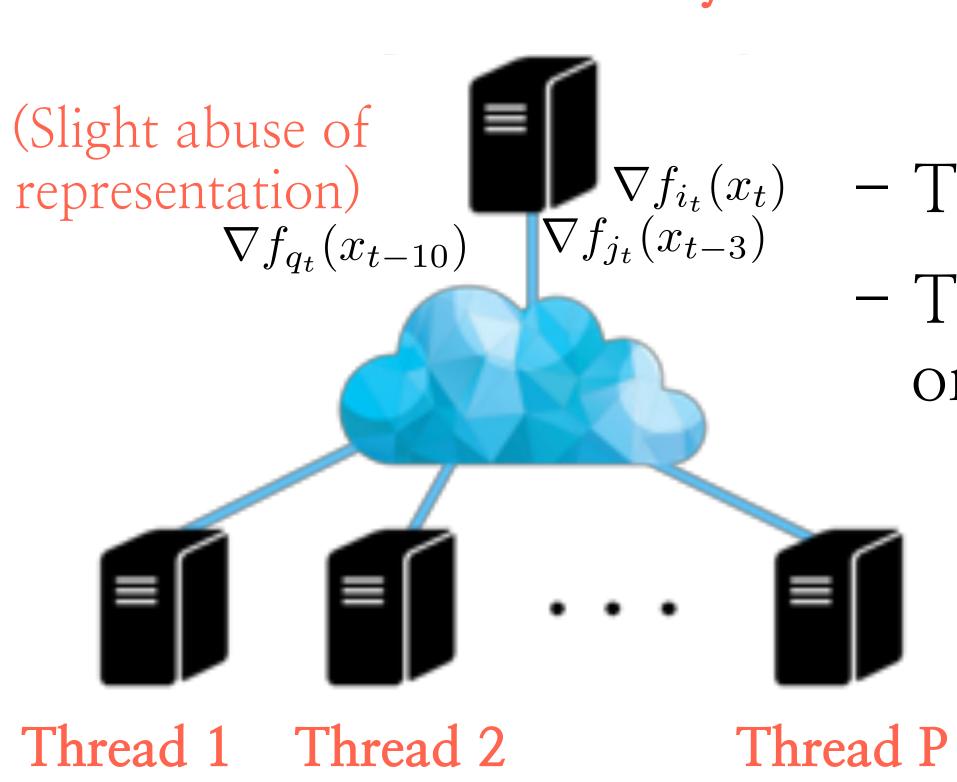


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- Threads might process an older model version

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- And it can get more complex:

"Threads can read a model state that only stayed in memory for a short time and between other memory writes"

(..and this might be an straightforward case)

Each has access to all data

– Does it work?

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(...a bit more involved set up)

#### Large Scale Distributed Deep Networks

Jeffrey Dean, Greg S. Corrado, Rajat Monga, Kai Chen, Matthieu Devin, Quoc V. Le, Mark Z. Mao, Marc'Aurelio Ranzato, Andrew Senior, Paul Tucker, Ke Yang, Andrew Y. Ng

{jeff, gcorrado}@google.com
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#### – Does it work?

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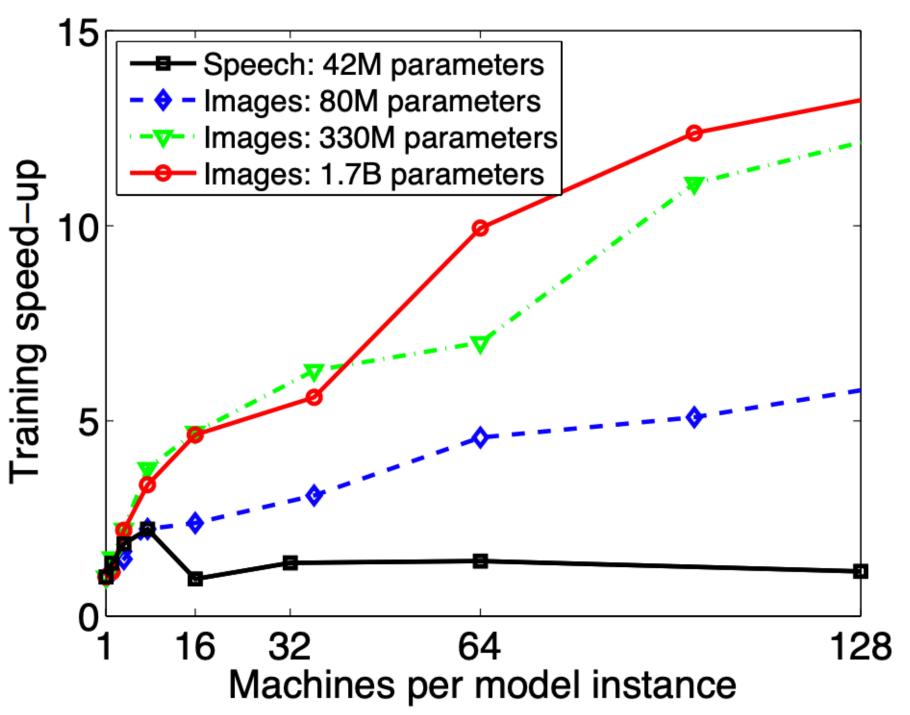
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- Comm. bottleneck
- We can increase batch
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- Setting: 
$$\min_{x} f(x) := \sum_{e \in E} f_e(x_e)$$
 where:  $x \in \mathbb{R}^n$   $E$  is a collection of items, say samples  $e \subset [n]$  (each element e is a collection of indices in [n] but also an index from a set of samples E)

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- Slight abuse of notation:

 $f_e(\cdot)$ : denotes a component of sum of functions, indexed by sample e

 $x_e$ : corresponds to sub-vector, indexed by an index set e (connected to sample e)

– Key observation: n & |E| are large, while individual  $f_e(\cdot)$  act on a small number of components of  $x \in \mathbb{R}^n$ 

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- Example: Sparse SVM

Given data  $E = \{(z_1, y_1), \dots, (z_{|E|}, y_{|E|}\}$  where  $y_i$  labels and  $z_i \in \mathbb{R}^n$  are features, we solve:

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– Observe that, if  $z_{\alpha}$  is very sparse (which happens in reality often), then

#### - Some quantities:

 $\Omega$ : maximum number of features involved over all samples

 $\Delta$ : maximum frequency of features that can appear in samples

 $\rho$ : approaches 1 if features are very common across examples

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### - Configuration:

p: number of processors

Each processor can read model x and contribute an update to x

#### Algorithm 1 Hogwild update for individual processors

- 1: **loop**
- 2: Sample e uniformly at random from E
- 3: Read current state  $x_e$  and evaluate  $G_e(x)$
- 4: for  $v \in e$  do  $x_v \leftarrow x_v \gamma b_v^T G_e(x)$  (coordinate-wise)
- 5: end loop

#### - Notation:

 $G_e(x) \in \mathbb{R}^n$ : gradient with non-zeros indexed by e, and scaled such that

$$\mathbb{E}\left[G_e(x_e)\right] = \nabla f(x)$$

Observe that  $[G_e(x_e)]_{e^c} = 0$ 

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#### - In words:

- 1. Each processor samples e uniformly at random
- 2. Each processor computes the gradient  $f_e$  at  $x_e$
- 3. Each processor applies update on each coordinate in e

– Asynchrony:  $x_j$  denotes the variable after j updates. Generally updated with stale gradients

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Whiteboard

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No Demo (no resources)

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  - The authors show (theoretically and experimentally) a near-linear speedup, with the number of processors used
  - In practice, lock-free SGD exceeds even theoretical guarantees

(in other words, how we can decrease communication burden?)

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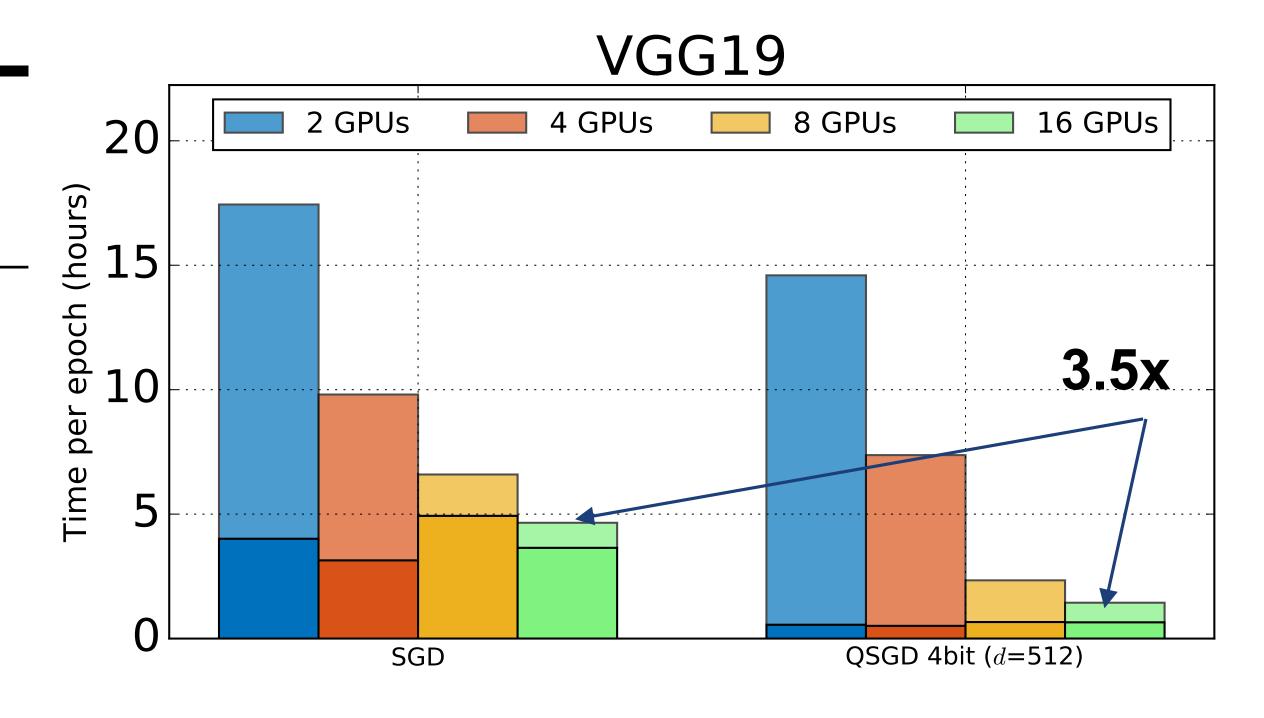
### QSGD: Communication-Efficient SGD via Gradient Quantization and Encoding

Dan Alistarh
IST Austria & ETH Zurich
dan.alistarh@ist.ac.at

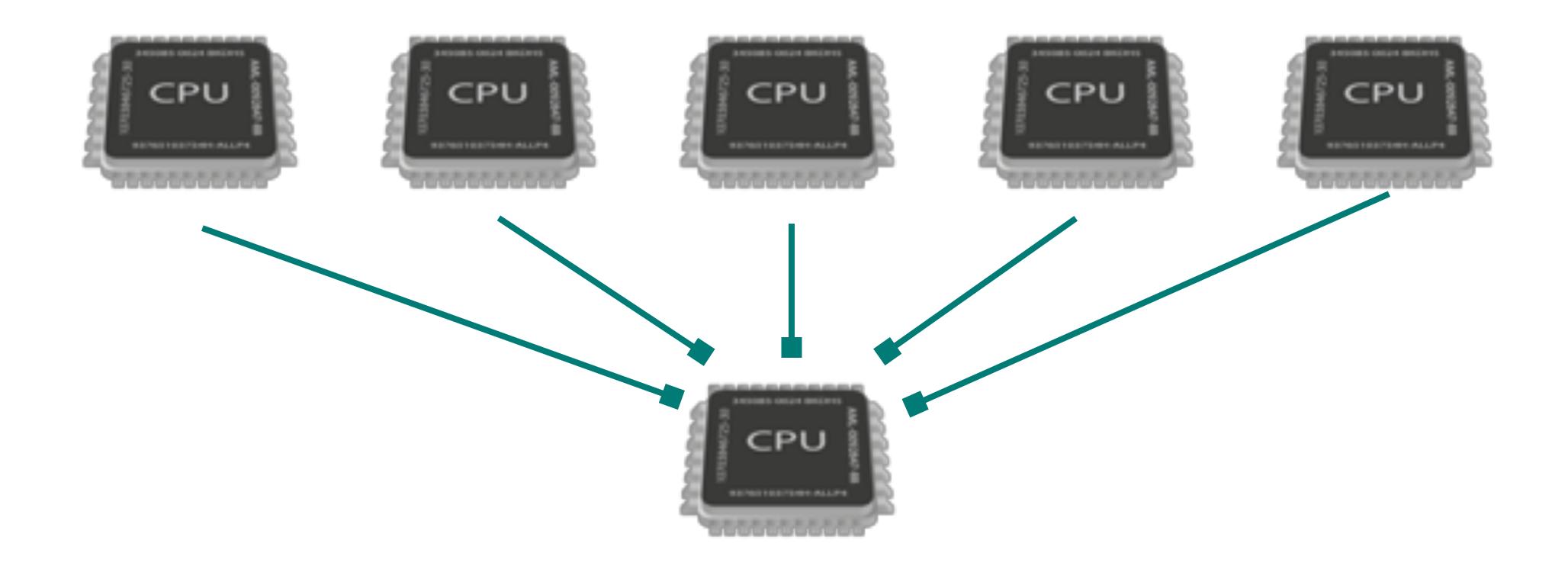
Demjan Grubic ETH Zurich & Google demjangrubic@gmail.com Jerry Z. Li
MIT
jerryzli@mit.edu

Ryota Tomioka
Microsoft Research
ryoto@microsoft.com

Milan Vojnovic
London School of Economics
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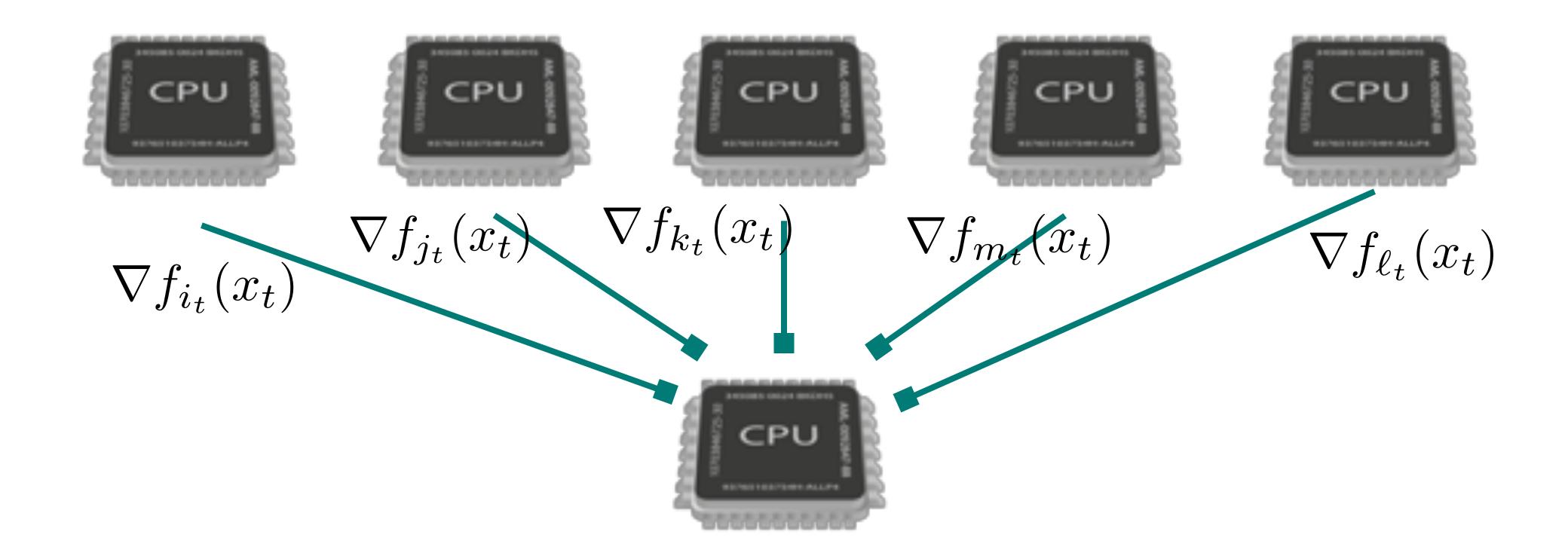


(in other words, can we make synchronization not be a big problem?)



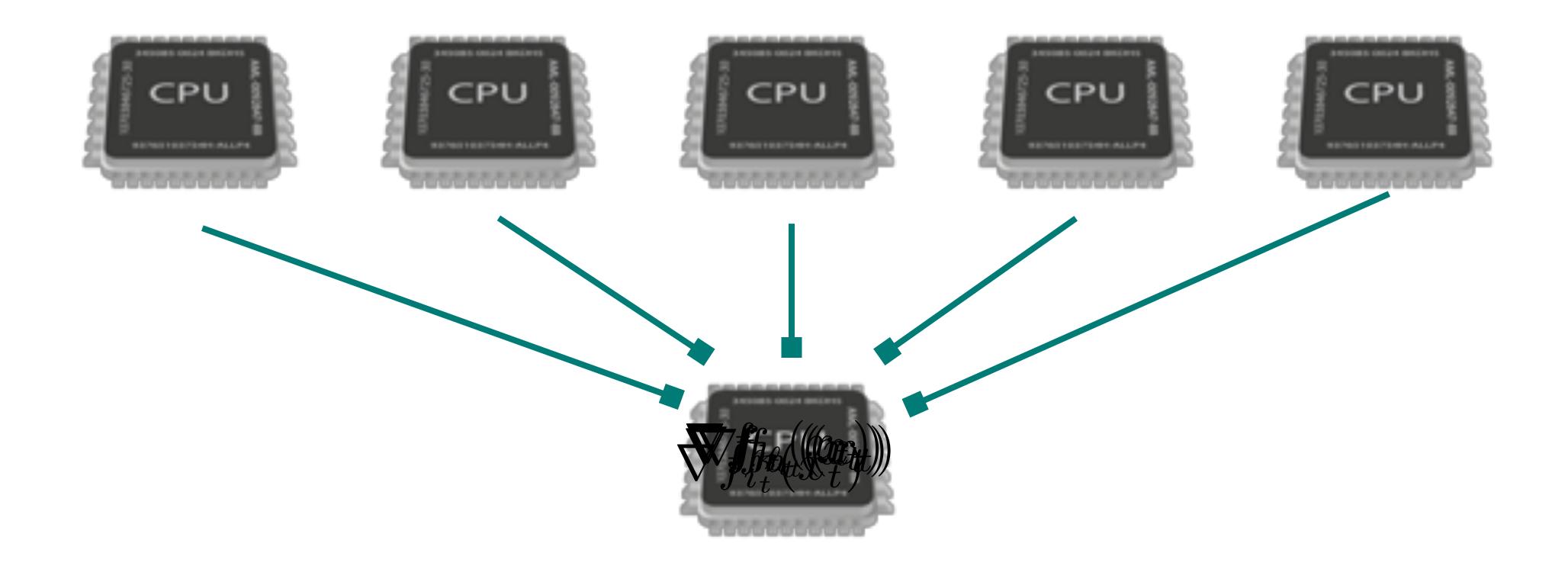
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#### REVISITING DISTRIBUTED SYNCHRONOUS SGD

Jianmin Chen, Xinghao Pan, Rajat Monga, Samy Bengio Google Brain
Mountain View, CA, USA
{jmchen, xinghao, rajatmonga, bengio}@google.com

Rafal Jozefowicz
OpenAI
San Francisco, CA, USA
rafal@openai.com

#### **ABSTRACT**

Distributed training of deep learning models on large-scale training data is typically conducted with *asynchronous* stochastic optimization to maximize the rate of updates, at the cost of additional noise introduced from asynchrony. In contrast, the *synchronous* approach is often thought to be impractical due to idle time wasted on waiting for straggling workers. We revisit these conventional beliefs in this paper, and examine the weaknesses of both approaches. We demonstrate that a third approach, synchronous optimization with backup workers, can avoid asynchronous noise while mitigating for the worst stragglers. Our approach is empirically validated and shown to converge *faster* and to *better* test accuracies.

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- Sparsification of gradients: instead of quantizing all entries, keep the most important ones
- Large batch training: give more "work" to workers by increasing the batch size. However it needs careful parameter tuning to make it work
- Variants of HOGWILD! that minimize communication conflicts: some computation is performed to distribute examples to different cores so that examples do not "conflict".

#### Conclusion

- Distributed computing is at the heart of developments in modern ML
- There are different ways to exploit distributed computing: hyper parameter optimization, coordinate descent, mini-batch synchronous SGD, asynchronous SGD
- Which configuration to use depends on the problem and the resources at hand
- These topics are highly attractive (research-wise): they define the notion of systems + machine learning (look for SysML conference)